

CSI 445/660 – Part 4  
(Positive and Negative Relationships)

Ref: Chapter 5 of [Easley & Kleinberg].

# Positive and Negative Relationships

- So far: Edges in a network represent **friendship** information (positive relationships).
- We also need to consider **conflicts** (negative relationships).
- The combination leads to the notion of **structural balance**.
- Provides another illustration of how **local** structure (i.e., a property involving a few nodes at a time) may have a **global** effect.

## Model:

- The underlying graph is a **clique**; that is, each person has a positive or negative relationship with every other person. (General graphs will be considered later.)
- Each edge has a **label**: '+' (indicating a positive relationship) or '-' (indicating a negative relationship).
- A common model for studying international conflicts.

## Model (continued):

- Ideas developed (in the sociological context) by Fritz Heider.



- Fritz Heider (1896–1988)
  - Austrian Sociologist.
  - Taught at the University of Kansas for many years.
- 
- The mathematical development is due to Dorwin Cartwright and Frank Harary.

## Positive and Negative ... (continued)

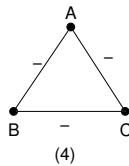
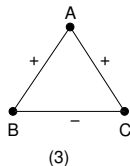
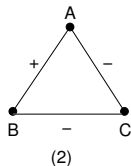
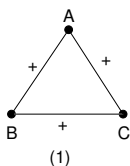


- Dorwin Cartwright (1915–2008)
- Areas: Psychology and Mathematics.
- One of the founders of Group Dynamics.
- University of Michigan, Ann Arbor, MI.



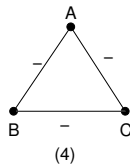
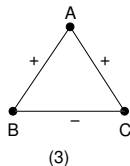
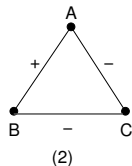
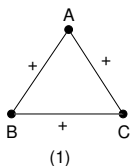
- Frank Harary (1921–2005)
- Mathematician who specialized in Graph Theory and its Applications.
- University of Michigan, Ann Arbor, MI and later New Mexico State University, Las Cruces, NM.

## Possible Edge Labelings for Three People:



- Labelings (1) and (2) have an **odd** number of '+' labels.
- Labelings (3) and (4) have an **even** number of '+' labels.
- Labeling (1): Three mutual friends; causes no problem.
- Labeling (2): Two friends and they both dislike the third; causes no problem.
- So, Labelings (1) and (2) have **structural balance**.

## Structural Balance (continued)



- Labeling (3): A has two friends who don't like each other. This may be a source of "stress" for A. (It may cause A to lose the friendship with B or C.)

**Note:** Recall (from the slides for Part 2) the study by Bearman & Moody [2004] involving the health records of teenage girls.

- Labeling (4): Here, two of the people may "team up" against the third person (i.e., there may be forces to change the label of one of the edges to '+').
- So, Labelings (3) and (4) have **structural imbalance**.

# Structural Balance (continued)

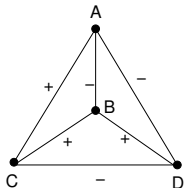
## Balance condition for Three People:

- A labeled triangle is **balanced** if and only if the number of '+' labels is **odd**.

## Extension – Structural Balance for all Cliques:

- A labeled clique is **balanced** if and only if each of its triangles is balanced (i.e., in each triangle, the number of '+' labels is **odd**).

## Example:



- 4-clique.
- Not balanced.
- Triangle BCD has two edges labeled '+' (and so does triangle ABC).

# Structural Balance (continued)

## Testing the Structural Balance – An Easy Algorithm:

**Input:** A clique  $G$  with  $n$  nodes where each edge has a '+' or '-' label.

**Output:** "Yes" if  $G$  is balanced and "No" otherwise.

### Outline of the Algorithm:

- 1** for each triple of nodes  $x, y$  and  $z$  **do**  
    **if** (triangle  $\{x, y, z\}$  is **not** balanced)  
        Output "No" and **stop**.
- 2** Output "Yes".

**Running time:**  $O(n^3)$  (since there are  $\binom{n}{3} = O(n^3)$  triangles in a clique with  $n$  nodes).



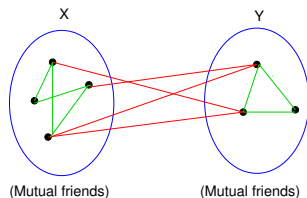
# Characterizing Structural Balance

**Note:** The following trivial cases are ignored in the discussion.

- All edges of  $G$  are labeled '+':  $G$  is balanced.
- All edges of  $G$  are labeled '-':  $G$  is unbalanced.

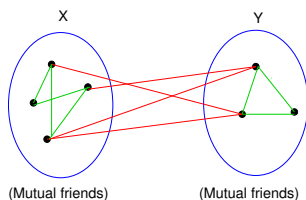
## Idea of Battling Factions:

- Suppose we can partition the nodes of  $G$  into two sets  $X$  and  $Y$  such that the following conditions hold:
  - All edges inside  $X$  or inside  $Y$  are labeled '+' and
  - all edges that join a node in  $X$  to a node in  $Y$  are labeled '-'.



- Not all edges are shown.
- Each **green** edge has the label '+', and each **red** edge has the label '-'.

# Characterizing Structural Balance (continued)



- Not all edges are shown.

- $X$  and  $Y$  are called **battling factions**.
- In this structure, every triangle is balanced (to be explained in class).
- **Key idea:** In any balanced clique, such a structure exists.

## Terminology:

- **Internal edge:** An edge that joins two nodes in  $X$  or two nodes in  $Y$ .
- **External edge:** An edge that joins a node in  $X$  to a node in  $Y$ .

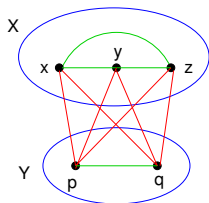
# Characterizing Structural Balance (continued)

## Theorem: [Cartwright & Harary]

If a labeled complete graph  $G$  is balanced, then

- either all the edge labels in  $G$  are '+' or
- the nodes of  $G$  can be partitioned into two sets  $X$  and  $Y$  such that
  - 1 each **internal** edge is labeled '+' and
  - 2 each **external** edge is labeled '-'.

## Example:



- This 5-clique is balanced.
- Partition:  $X = \{x, y, z\}$  and  $Y = \{p, q\}$ .

# Proof Sketch for the Cartwright-Harary Theorem

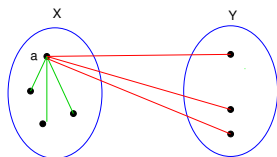
## Notes:

- Ignore the (trivial) case where all edge labels are '+'.  
■ So, assume that at least one edge has the label '-'.  
■ The proof actually constructs the **battling factions** partition.

## Construction:

- Choose any node  $a$  of  $G$ .  
■ Let the set  $X$  consist of  $a$  and all the nodes which are **friends** of  $a$ .  
■ Let  $Y$  be the remaining set of nodes (i.e., the **enemies** of  $a$ ).

## An Illustration:

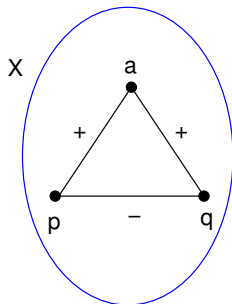


- Not all nodes/edges are shown.

# Proof Sketch ... (continued)

**Part 1:** We must show that each internal edge in  $X$  has the label '+'.  
■

- Consider any two nodes  $p$  and  $q$  in  $X$ .
- If one of  $p$  and  $q$  is the node  $a$ , the conclusion follows since all nodes in  $X$  are friends of  $a$ .
- So, assume that  $p$  and  $q$  are different from  $a$ .
- If  $p$  and  $q$  are enemies, we get the following **unbalanced** triangle in  $G$ :

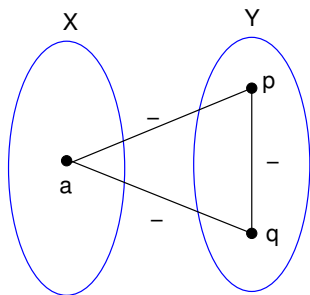


- This contradicts the assumption that  $G$  is balanced.

# Proof Sketch ... (continued)

**Part 2:** We must show that each internal edge in  $Y$  has the label '+'.  
■ Consider any two nodes  $p$  and  $q$  in  $Y$ .

- If  $p$  and  $q$  are enemies, we get the following **unbalanced** triangle in  $G$ :

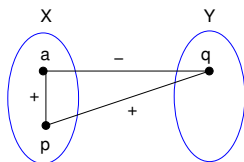


- A contradiction.

# Proof Sketch ... (continued)

**Part 3:** We must show that each external edge has the label '-'.

- Consider any two nodes  $p \in X$  and  $q \in Y$ .
- If  $p$  and  $q$  are friends, we get the following **unbalanced** triangle in  $G$ :



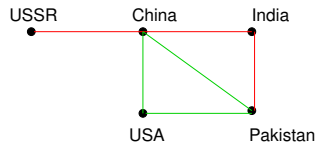
- A contradiction. (This completes the proof.)

## Notes:

- The Cartwright-Harary Theorem leads to an  $O(n^2)$  algorithm for the problem. (See Handout 4.1.)
- The running time is **linear** in the **size** of the input.

# An Application – International Relations

Ref: **[Moore 1978]** (Reference [308] in the text).



- Relationships in 1972.
- There was a war between India and Pakistan.
- USA was trying to improve its relationship with China.
- The perception was that China and Pakistan were friends (since India was their common 'enemy').
- The structural balance theory suggests that USA should support Pakistan.



## Other Related Topics (Brief Discussion)

- Some online networks allow people to express positive/negative sentiments.

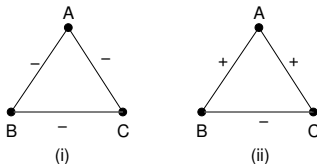
### Examples:

- Slashdot (<http://slashdot.org>): Allows people to designate each other as 'friend' or 'foe'.
- Epinions (<http://www.epinions.com>): A consumer review website where people could 'trust' or 'distrust' reviews. (These features were removed in 2014.)

- Evolving models of signed graphs  
(e.g. [Antal et al. 2006] – Ref [20] in the text).
  - 1 Start with a random labeling.
  - 2 Look for an unbalanced triangle and flip one of the labels to make it balanced.
  - 3 Repeat Step 2 until all triangles are balanced (or until the number of repetitions exceeds a set limit).
- Capture situations where people update their likes/dislikes as they strive for structural balance.

# A Weaker Form of Structural Balance

- Two forms of structural **imbalance**:



- Some sociologists (e.g. James Davis, University of Chicago) have argued that (ii) is a stronger form of imbalance than (i).

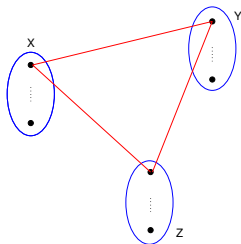
## Definition: [Weaker form of Imbalance]

A clique with signed edges is **weakly balanced** if and only if there is no triangle with **exactly two** edges labeled '+'.

**Note:** One should expect a larger collection of possible structures that are weakly balanced.

# A Weaker Form of Structural Balance (continued)

## Example:



- This structure **allows** triangles with three edges labeled '-'.
- However, triangles with only one edge labeled '-' are **not allowed**.

## Characterization of Weakly Balanced Cliques:

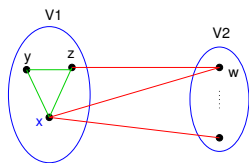
### Theorem: (Also due to Cartwright & Harary)

Let  $G$  be a weakly balanced clique. Then the nodes of  $G$  **can be partitioned into groups** such that for any pair of nodes  $x$  and  $y$

- 1 if  $x$  and  $y$  are in the same group, then  $x$  and  $y$  are friends and
- 2 if  $x$  and  $y$  are in different groups, then  $x$  and  $y$  are enemies.

# Weak Balance for Cliques: Proof Idea

- Let  $G$  be the weakly balanced clique.
- Choose any node  $x$  of  $G$  and construct set  $V1$  consisting of  $x$  and all the friends of  $X$ .
- Let  $V2$  denote the set of remaining nodes.



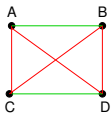
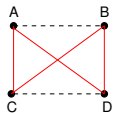
- Each pair of nodes in  $V1$  must be friends. (Otherwise, will have a triangle in  $V1$  with exactly one edge labeled '-', which is **not** weakly balanced.)
  - Also, for any node  $z \in V1$  and any node  $w \in V2$ ,  $z$  and  $w$  are enemies.
- Think of  $V1$  as the first group.
  - The complete graph on  $V2$  is also **weakly balanced**. So, one can continue the process with  $V2$ , leading to several groups.

# Strong Structural Balance for General Graphs

- **So far:** Balance conditions for cliques.
- **Now:** Strong structural balance for graphs which are not necessarily cliques.
- There are two possible definitions.

**Definition 1:** Let  $G$  a graph with each edge labeled '+' or '-'.  $G$  is **balanced** if signs can be assigned to the **missing edges** so that the resulting clique is (strongly) balanced.

**Example:**



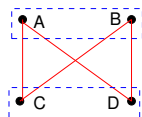
- The graph on the right assigns the '+' label to each missing edge.
- So, the graph on the left is **balanced**.

## Balance for General Graphs (continued)

**Definition 2:** Let  $G$  a graph with each edge labeled '+' or '-'.  $G$  is **balanced** if the nodes of  $G$  can be partitioned into two sets  $V_1$  and  $V_2$  such that

- 1 Each edge inside  $V_1$  or  $V_2$  has the '+' label and
- 2 each edge that joins a node in  $V_1$  to a node in  $V_2$  has the '-' label.

**Example:**



**Note:** There need not be any internal edges.

**Fact:** Definitions 1 and 2 are **equivalent**; that is, a graph  $G$  is balanced according to Definition 1 and if and only if it is balanced according to Definition 2.

## Reason for the Equivalence of Definitions:

- If it is possible to assign labels to missing edges so that the graph becomes balanced (by Definition 1), then we can obtain a “battling factions” partition that satisfies Definition 2.
- If the graph satisfies Definition 2, then all internal edges can be labeled ‘+’ and all external edges can be labeled ‘-’ to satisfy Definition 1.

**Note:** Unfortunately, these definitions don’t directly lead to an efficient algorithm for checking the balance condition for general graphs.

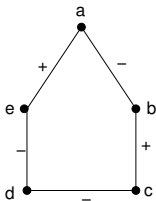
## Theorem: [Harary]

A signed graph is balanced if and only if it does **not** contain any cycle with an **odd** number of edges with label ‘-’.



# Explanation for Harary's Theorem

**Example:** The following graph has a cycle with an **odd** number of edges labeled '-':



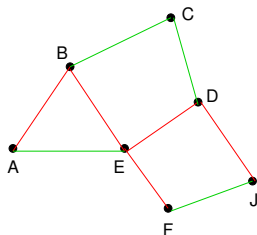
- In any “battling factions” decomposition, nodes *a* and *e* must be on the same side.
- Likewise, nodes *b* and *c* must be on the same side, but **different** from the side that contains *a* and *e*.
- Now, we can't add node *d* to either side.
- So, the above graph is **not** balanced.

**Note:** Harary's theorem leads to an efficient algorithm for testing the strong balance condition for general graphs.

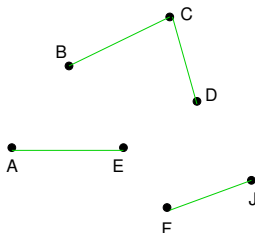
**Algorithm Description:** See Handout 4.2.

# An Illustration for the Algorithm

Given signed graph  $G$ :

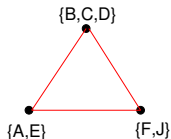


Graph  $G^+$  after Step 2:



**Note:**  $G$  does not contain any edge labeled '-' joining two nodes in the same connected component.

Graph  $H$  after Step 4:



- $H$  is **not** bipartite; it contains a cycle with 3 nodes.
- So,  $G$  is not balanced.

# Notion of Approximate Balance (Brief Discussion)

- **So far:** “Perfect balance” (i.e., **all** triangles are balanced).
- Suppose we allow 0.1% of “unbalanced” triangles; that is, in the given signed clique  $G$ , 99.9% of the triangles are balanced. Then, the following result holds.

**Theorem:** Suppose  $G$  is a signed clique such that 99.9% of the triangles in  $G$  satisfy the strong balance condition. Then **at least one** of the following conditions hold:

- There is a subset  $V'$  with at least 90% of the nodes of  $G$  such that at least 90% of the edges inside  $|V'|$  are labeled ‘+’.
- The nodes of  $G$  can be partitioned into two sets  $V_1$  and  $V_2$  such that
  - 1 at least 90% of the internal edges are labeled ‘+’ and
  - 2 at least 90% of the external edges are labeled ‘-’.

**Note:** A proof of the above result is given in the text.