Cyber-Physical Systems

Discrete Dynamics

IECE 553/453– Fall 2019

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Discrete Systems

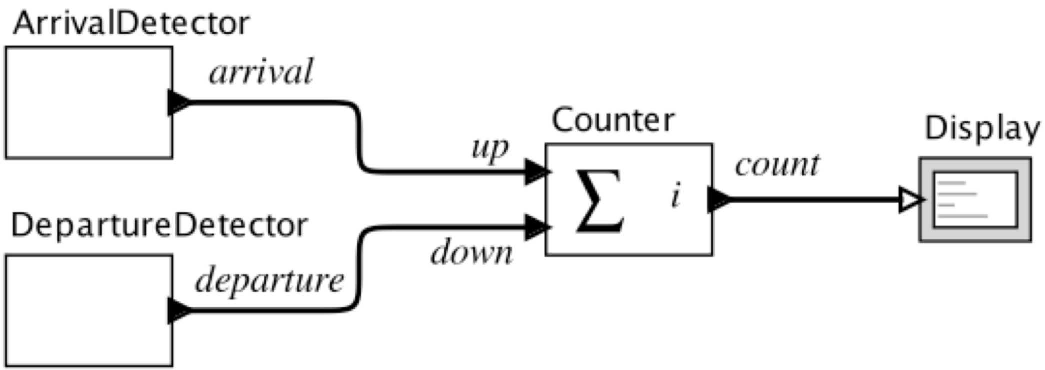
- **Discrete** = "individually separate / distinct"
- A **discrete system** is one that operates in a sequence of discrete *steps* or has signals taking discrete *values*.
- It is said to have **discrete dynamics**.
- A discrete event occurs at an instant of time rather than over time.

Discrete Systems: Example Design Problem

- Count the number of cars that are present in a parking garage by sensing cars enter and leave the garage. Show this count on a display.

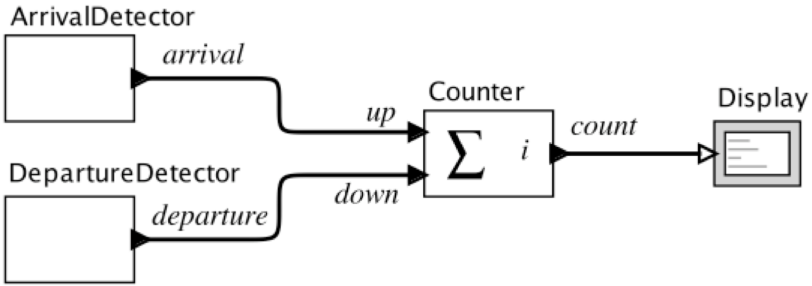
Discrete Systems

- Example: count the number of cars in a parking garage by sensing those that enter and leave:



Discrete Systems

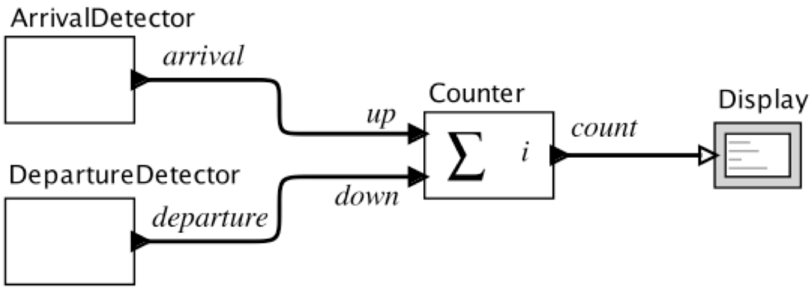
- Example: count the number of cars that enter and leave a parking garage:



- Pure signal: $up : \mathbb{R} \rightarrow \{absent, present\}$
- absent: no event at that time ; present: event at that time

Discrete Systems

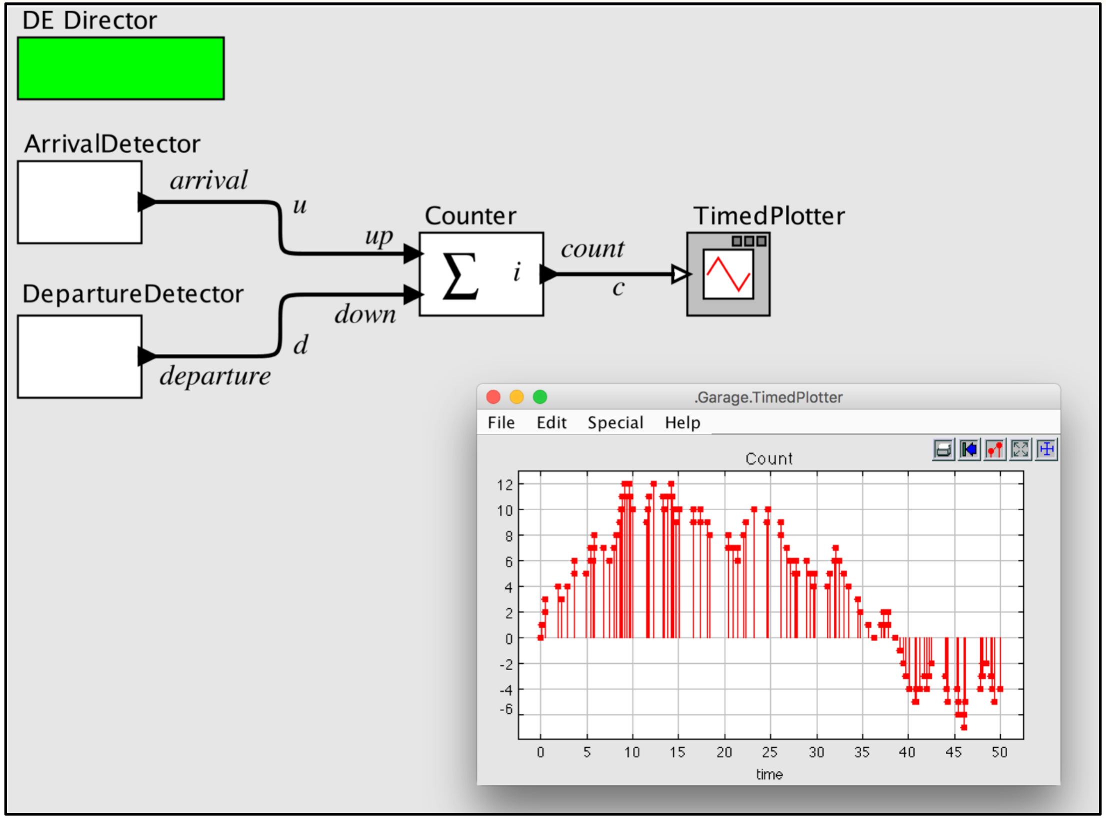
- Example: count the number of cars that enter and leave a parking garage:



- Pure signal: $up: \mathbb{R} \rightarrow \{absent, present\}$
- Discrete actor: $Counter: (\mathbb{R} \rightarrow \{absent, present\})^P \rightarrow (\mathbb{R} \rightarrow \{absent\} \cup \mathbb{N})$

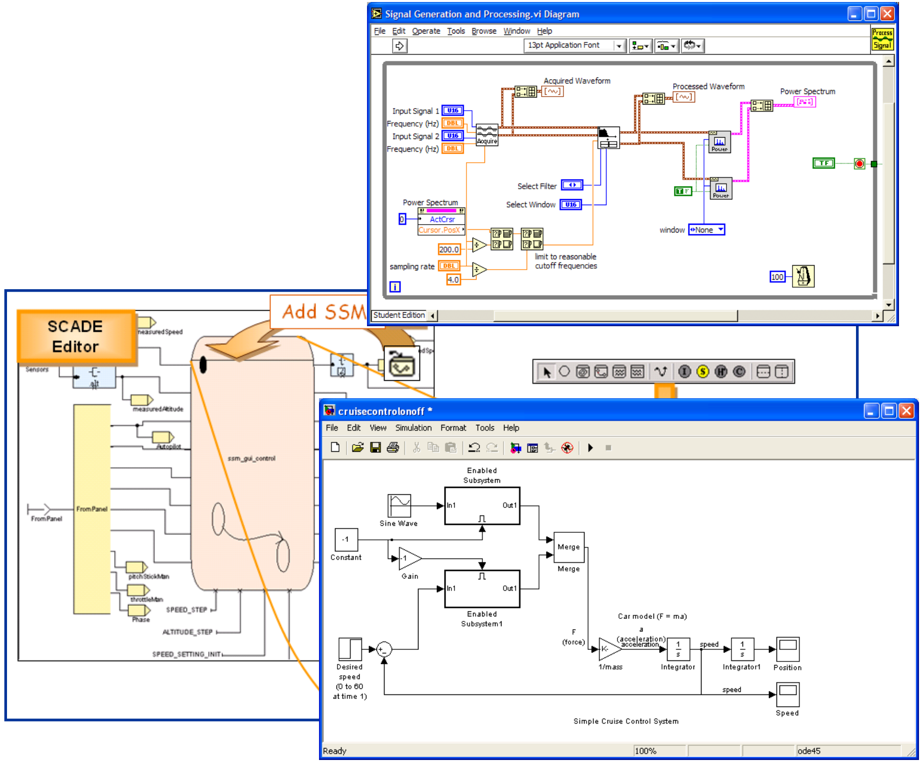
  $P = \{up, down\}$

Demonstration of Ptolemy II Model ("Program")



Actor Modeling Languages / Frameworks

- LabVIEW
- Simulink
- Scade
- ...
- Reactors
- StreamIT
- ...

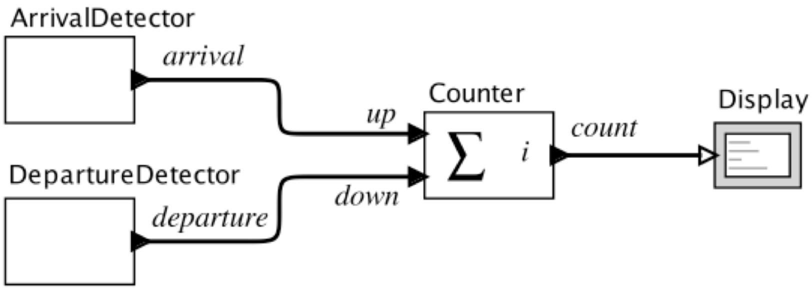


Reaction / Transition

For any $t \in \mathbb{R}$ where $up(t) \neq absent$ or $down(t) \neq absent$ the Counter **reacts**. It produces an output value in $\mathbb{N}$ and changes its internal **state**.

State: condition of the system at a particular point in time

- Encodes everything about the past that influences the system's reaction to current input



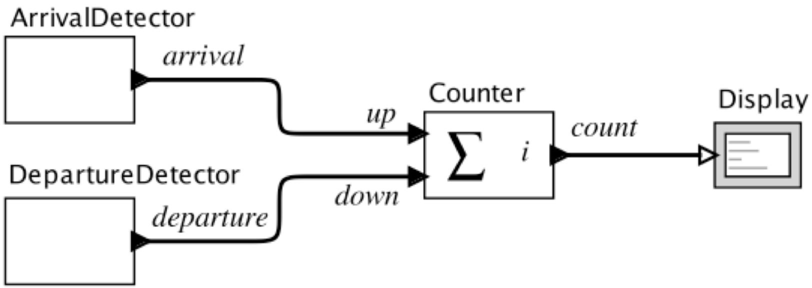
Inputs and Outputs at a Reaction

For $t \in \mathbb{R}$ the inputs are in a set

$$Inputs = (\{\mathsf{up}, \mathsf{down}\} \to \{\, absent, present\})$$

and the outputs are in a set

$$Outputs = (\{count\} \to \{\, absent\} \cup \mathbb{N})\ ,$$



Question

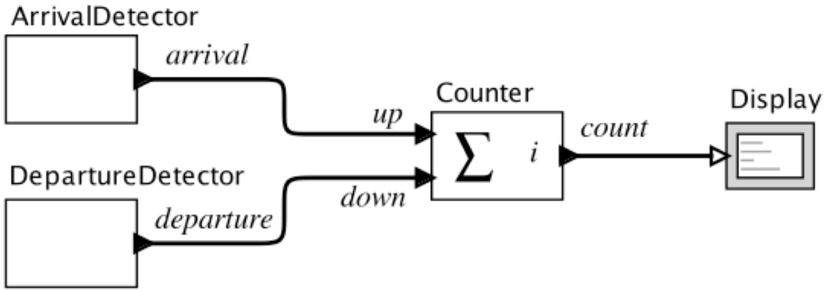
- What are some scenarios that the given parking garage (interface) design does not handle well?

For $t \in \mathbb{R}$ the inputs are in a set

$$Inputs = (\{\mathsf{up}, \mathsf{down}\} \rightarrow \{\, absent, present\})$$

and the outputs are in a set

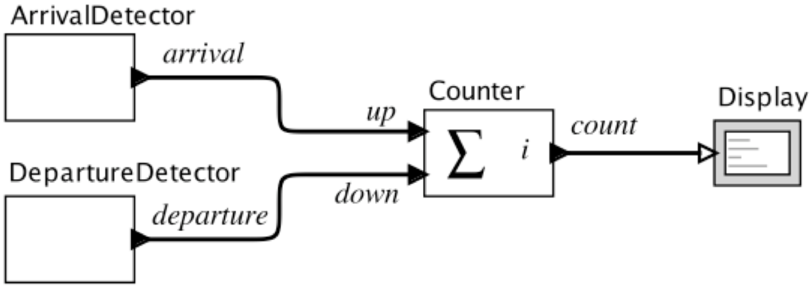
$$Outputs = (\{count\} \rightarrow \{\, absent\} \cup \mathbb{N}) \ ,$$



State Space

A practical parking garage has a finite number M of spaces, so the state space for the counter is

$$States = \{0, 1, 2, \cdots, M\} \ .$$



Finite State Machine (FSM)

- A state machine is a model of a system with discrete dynamics
  - at each reaction maps inputs to outputs
  - Map may depend on current state
- An FSM is a state machine where the set *States* is finite.

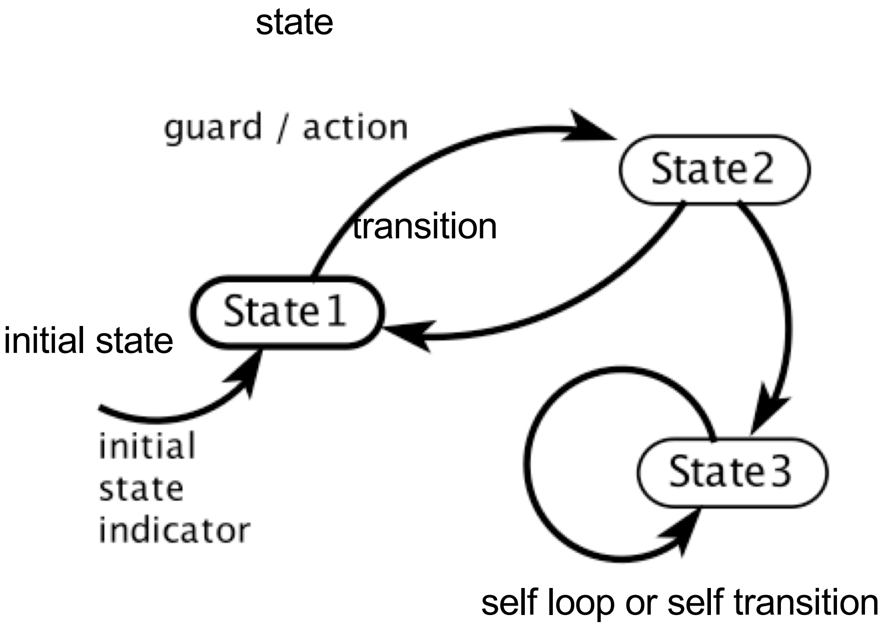
$$States = \{\text{State1}, \text{State2}, \text{State3}\}$$

FSM Notation

Input declarations, Output declarations, Extended state declarations

The guard determines whether the transition may be taken on a reaction.

The action specifies what outputs are produced on each reaction.



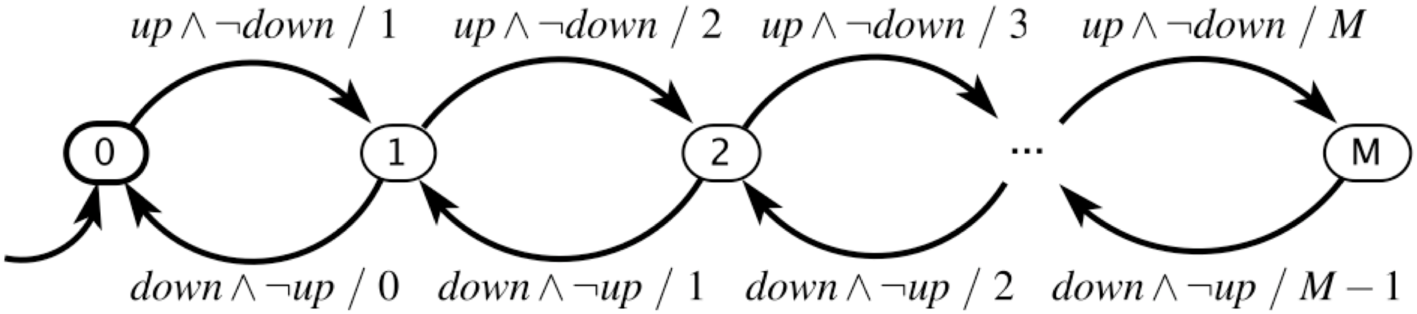
Examples of Guards for Pure Signals

true	Transition is always enabled.
p_1	Transition is enabled if p_1 is $present$.
$\neg p_1$	Transition is enabled if p_1 is $absent$.
$p_1 \wedge p_2$	Transition is enabled if both p_1 and p_2 are $present$.
$p_1 \vee p_2$	Transition is enabled if either p_1 or p_2 is $present$.
$p_1 \wedge \neg p_2$	Transition is enabled if p_1 is $present$ and p_2 is $absent$.

Guards for Signals with Numerical Values

p_3	Transition is enabled if p_3 is *present* (not *absent*).
$p_3 = 1$	Transition is enabled if p_3 is *present* and has value 1.
$p_3 = 1 \land p_1$	Transition is enabled if p_3 has value 1 and p_1 is *present*.
$p_3 > 5$	Transition is enabled if p_3 is *present* with value greater than 5.

Garage Counter Finite State Machine (FSM)



Guard $g \subseteq Inputs$ is specified using the shorthand

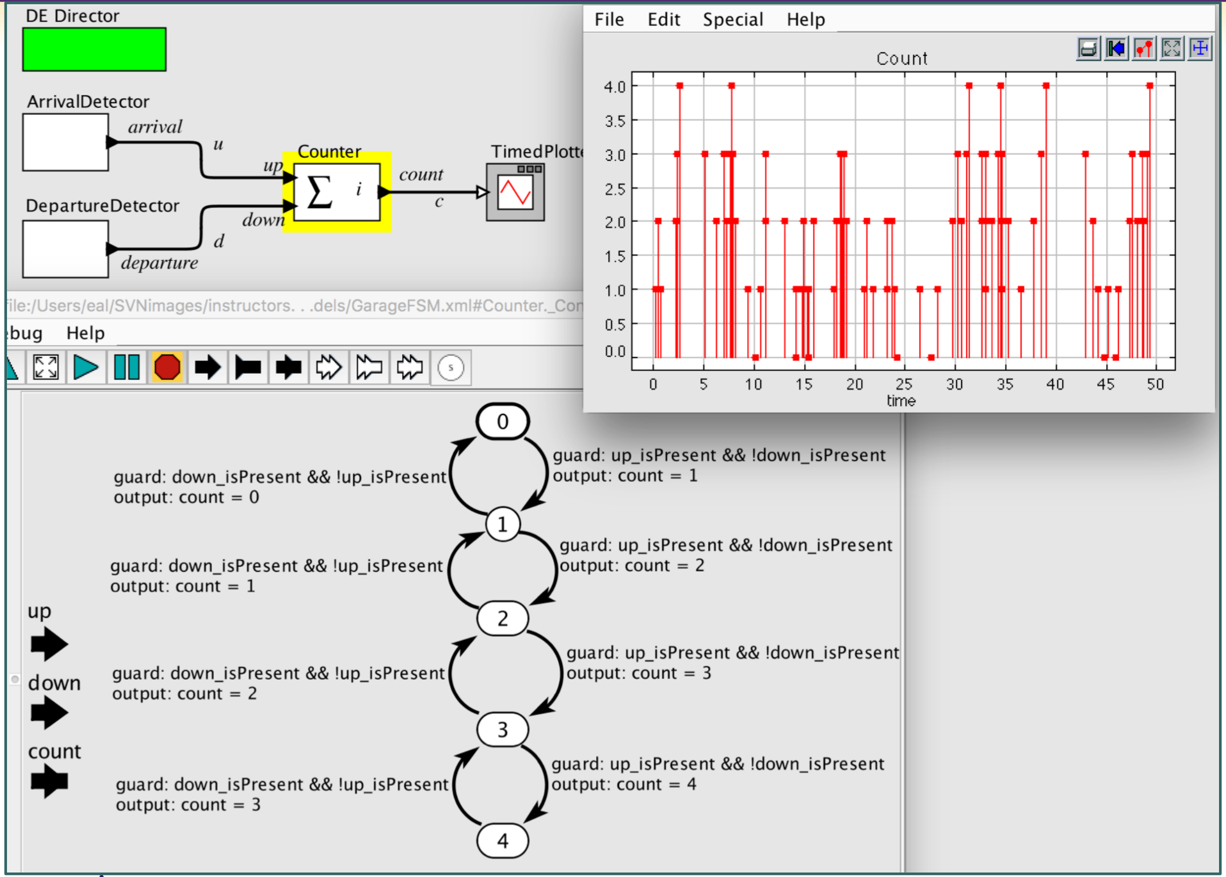
$$up \wedge \neg down$$

which means

$$g = \{\{up\}\} \ .$$

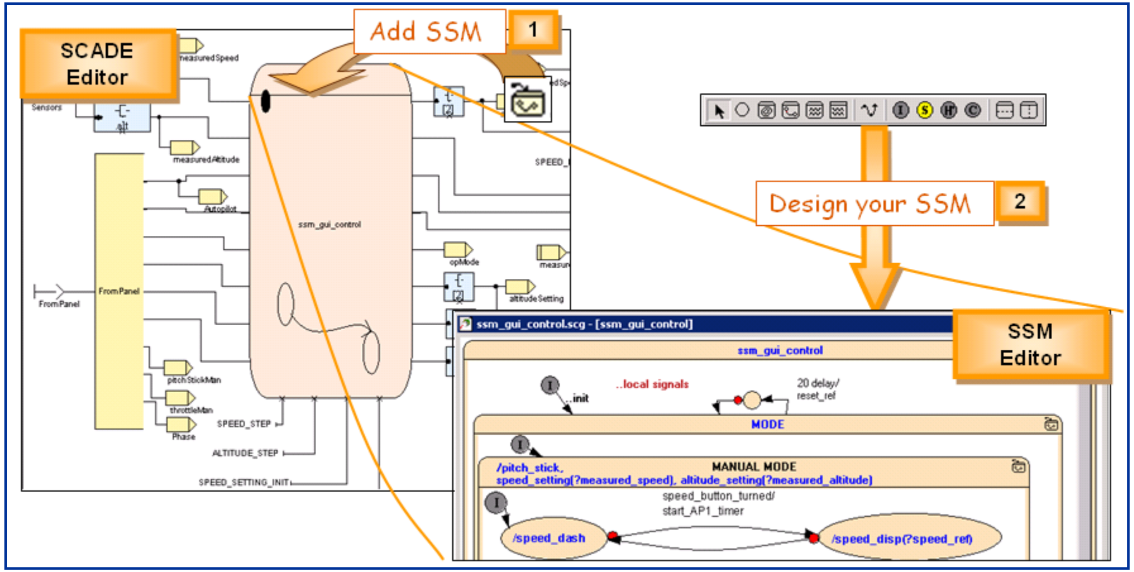
Inputs(up) = present and Inputs(down) = absent

Ptolemy II Model

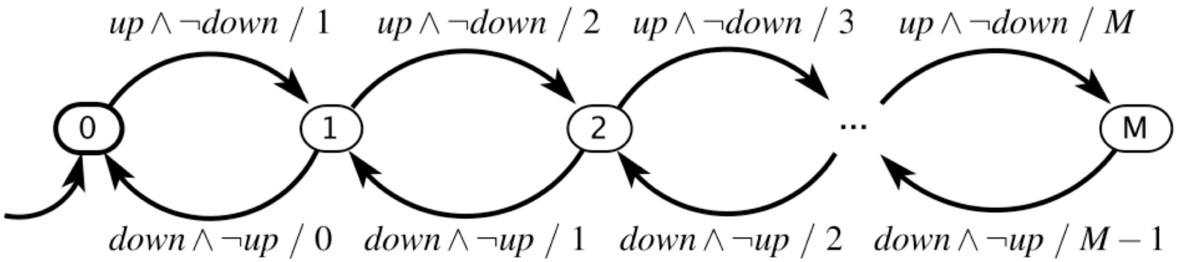


FSM Modeling Languages / Frameworks

- LabVIEW Statecharts
- Simulink Stateflow
- Scade
- …



Garage Counter Mathematical Model



Formally: $(States, Inputs, Outputs, update, initialState)$, where

- $States = \{0, 1, \cdots, M\}$
- $Inputs = (\{\text{up}, \text{down}\} \to \{absent, present\}$
- $Outputs = (\{count\} \to \{absent\} \cup \mathbb{N})$
- $update : States \times Inputs \to States \times Outputs$
- $initialState = 0$

Transition Function

$$\boxed{(s(n+1), y(n)) = update(s(n), x(n))}$$

The update function is given by

$$update(s,i) = \begin{cases} (s+1, s+1) & \text{if } s < M \\ & \land\, i(up) = present \\ & \land\, i(down) = absent \\ (s-1, s-1) & \text{if } s > 0 \\ & \land\, i(up) = absent \\ & \land\, i(down) = present \\ (s, absent) & \text{otherwise} \end{cases}$$

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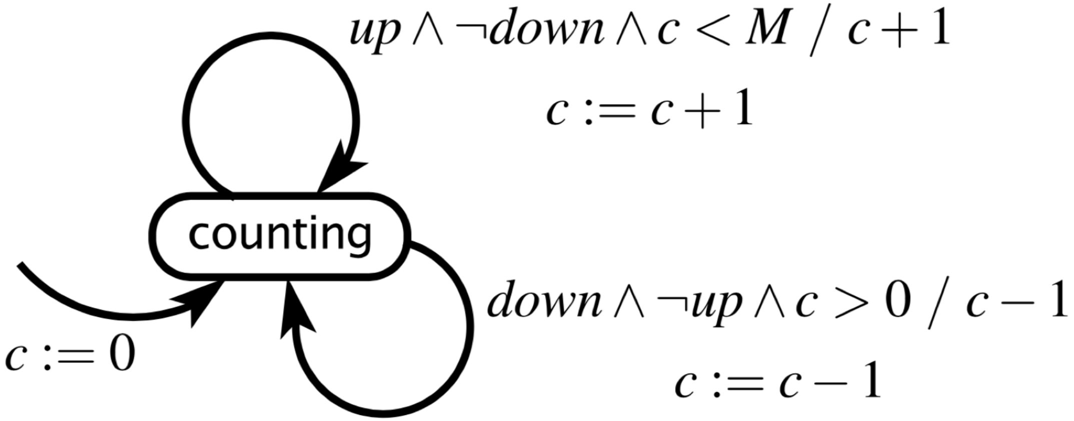
FSM: Determinacy and Receptiveness

- Deterministic (given the same inputs it will always produce the same outputs)
  - if, for each state, there is at most one transition enabled by each input value.
  - The formal definition of an FSM ensures that it is deterministic, since *update* is a function.
- Receptive (ensures that a state machine is always ready to react to any input, and does not "get stuck" in any state)
  - if, for each state, there is at least one transition possible on each input symbol.
  - The formal definition of an FSM ensures that it is receptive, since *update* is a function, not a partial function.

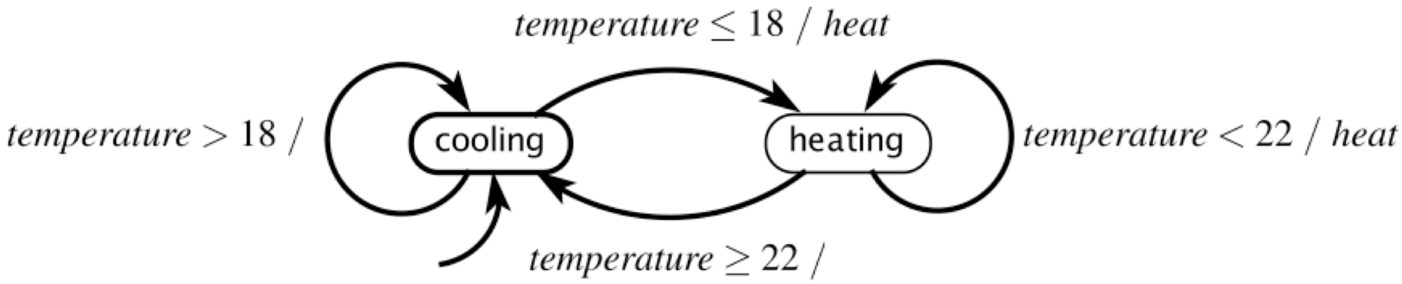
Extended State Machine

- augments the FSM model with variables that may be read and written as part of taking a transition between states

variable: $c : \{0, \cdots, M\}$
inputs: $up, down$: pure
output: $count : \{0, \cdots, M\}$



Example of Thermostat

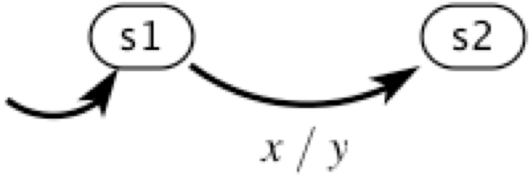


When does a reaction occur?

- Suppose all inputs are discrete and a reaction occurs *when any input is present*. Then the below transition will be taken whenever the current state is s1 and *x* is present.

- This is an *event-triggered model*.

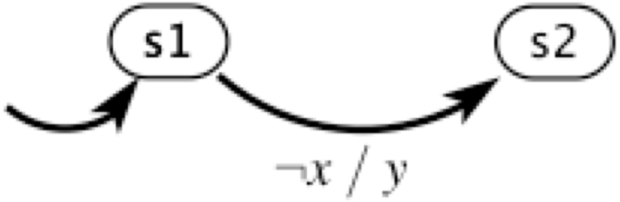
input: $x \in \{present, absent\}$
output: $y \in \{present, absent\}$



When does a reaction occur?

- Suppose *x* and *y* are discrete and pure signals. When does the transition occur?

input: $x \in \{present, absent\}$
output: $y \in \{present, absent\}$



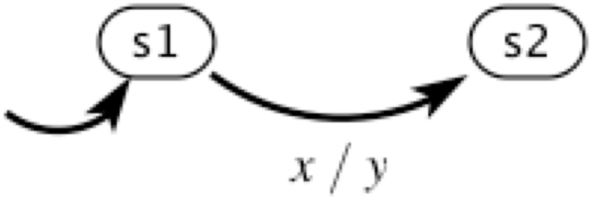
Answer: when the *environment* triggers a reaction and x is absent. If this is a (complete) event-triggered model, then the transition will never be taken because the reaction will only occur when x is present!

When does a reaction occur?

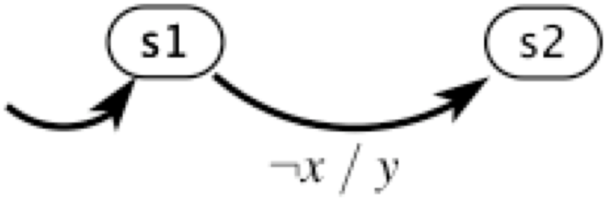
- Suppose all inputs are discrete and a reaction occurs *on the tick of an external clock.*

- This is a *time-triggered model.*

input: $x \in \{present, absent\}$
output: $y \in \{present, absent\}$

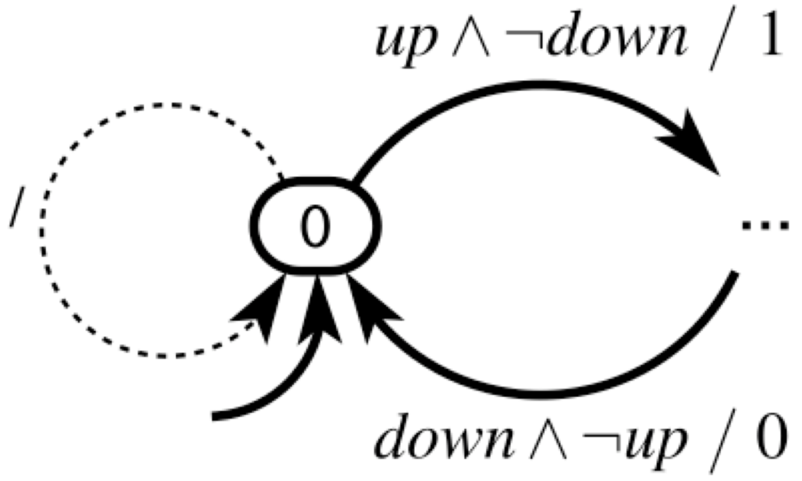


input: $x \in \{present, absent\}$
output: $y \in \{present, absent\}$



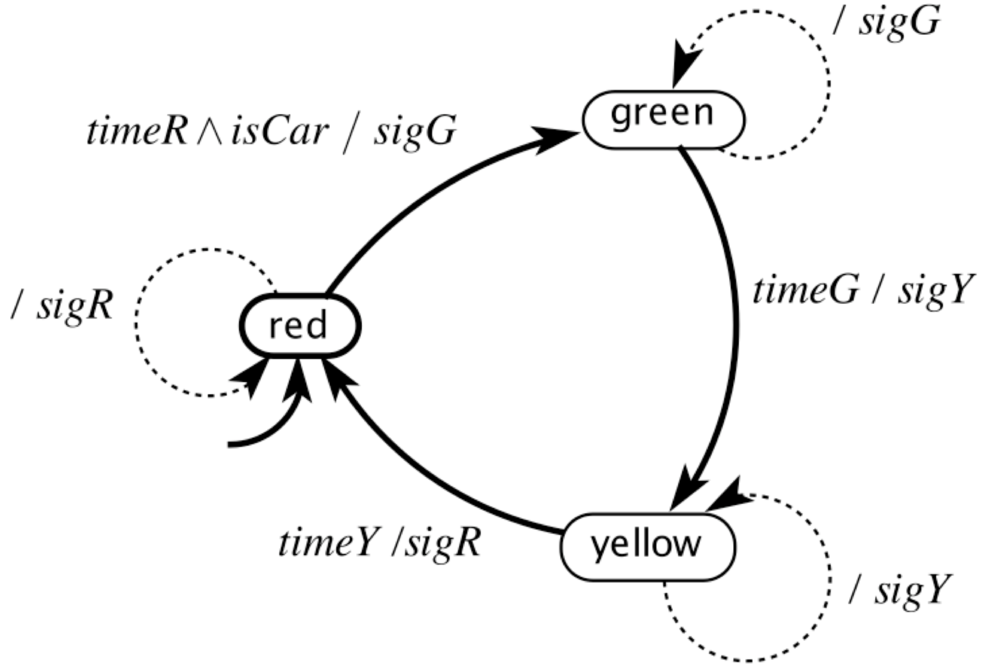
More Notation: Default Transitions

- A default transition is enabled if it either has no guard or the guard evaluates to true. When is the below default transition enabled?

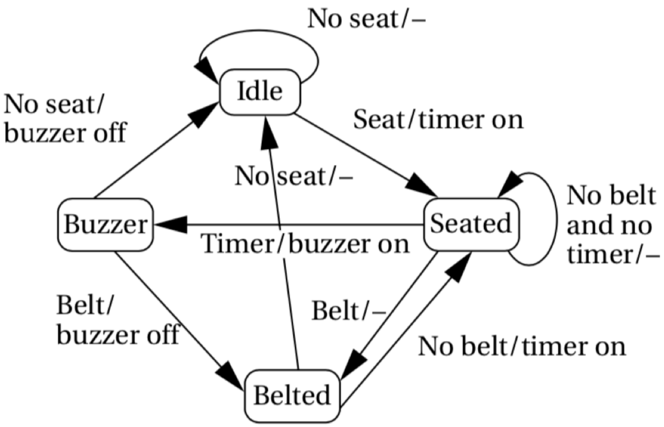


Default Transitions

- Example: Traffic Light Controller



FSM to Program




```c
#define IDLE 0
#define SEATED 1
#define BELTED 2
#define BUZZER 3

switch (state) { /* check the current state */
        case IDLE:
                if (seat) { state = SEATED; timer_on = TRUE; }
                /* default case is self-loop */
                break;
        case SEATED:
                if (belt) state = BELTED; /* won't hear the
                                 buzzer */
                else if (timer) state = BUZZER; /* didn't put on
                                         belt in time */
                /* default is self-loop */
                break;
        case BELTED:
                if (!seat) state = IDLE; /* person left */
                else if (!belt) state = SEATED; /* person still
                                         in seat */
                break;
        case BUZZER:
                if (belt) state = BELTED; /* belt is on—turn off
                                 buzzer */
                else if (!seat) state = IDLE; /* no one in
                                         seat—turn off buzzer */
                break;
}
```