# Combination Testing Strategies: A Survey * 

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#### Abstract

Combination strategies are test case selection methods that identify test cases by combining values of the different test object input parameters based on some combinatorial strategy. This survey presents 16 different combination strategies, covering more than 40 papers that focus on one or several combination strategies. This collection represents most of the existing work performed on combination strategies. This survey describes the basic algorithms used by the combination strategies. Some properties of combination strategies, including coverage criteria and theoretical bounds on the size of test suites, are also included in this description.

This survey paper also includes a subsumption hierarchy that attempts to relate the various coverage criteria associated with the identified combination strategies.


Keywords: Combination Strategies, Category Partition, Orthogonal Arrays, AETG, Test Case Selection

## 1 Introduction

Combination strategies define ways to select values for individual input parameters and combine them to form complete test cases. A literature search has revealed 16 different combination strategies described through more than 40 papers published between 1985 and 2004. This survey is an attempt to collect all this knowledge in one single place and form a basis for comparisons among the combination strategies.

[^0]Combination strategies are useful in many settings, not only in software testing. This survey only includes papers that explicitly mention software testing.

Section 2 gives a brief background on testing in general and on the Category Partition method [1] in particular. The Category Partition method is a convenient place to start because it yields independent values for parameters and spawned much of the work in combination strategies. Section 3 starts by explaining the coverage criteria usually associated with combination strategies, then gives an overview of each combination strategy identified in this survey. To aid the reader the combination strategies have been structured into different groups based on different properties of the combination strategy algorithms. The overview of each algorithm includes an explanation of the algorithm and an example. Associated coverage criteria are also included in the description. Section 4 compares the different combination strategies with respect to size of generated test suite. A new contribution of this paper is a subsumption hierarchy for the identified coverage criteria. It can be used to compare the combination strategies with respect to their associated coverage criteria.

Section 5 summarizes the collected experiences of using combination strategies. Finally, section 6 concludes this survey with a summary of the most important results and some future directions of research.

## 2 Background

Testing, loosely considered to be the dynamic execution of test cases, consumes a significant amount of the resources in software projects [2,3]. Thus, it is of great importance to investigate ways of increasing the efficiency and effectiveness in testing $[4,5,6,7,8]$.

Several existing testing methods (e.g. Equivalence Partitioning [2], Category Partition [1], and Domain Testing [3]) are based on the model that the input space of the test object may be divided into subsets based on the assumption that all points in the same subset result in a similar behavior from the test object. This is called partition testing. Even if the partition test assumption is an idealization it has two important properties. First, it lets the tester identify test suites of manageable size by selecting one or a few test cases from each subset. Second, it allows test effectiveness to be measured in terms of coverage with respect to the partition model used.

One alternative to partition testing is random testing, in which test cases are chosen randomly from some input distribution (such as a uniform distribution) without exploiting information from the specification or previously chosen test cases. Duran and Ntafos [9] gave results that under certain very restrictive conditions, random testing can be as effective as partition testing. They showed consistent small differences in effectiveness between partition testing methods and random testing. These results were interpreted in favor of random testing since it is generally less work to construct test cases in random testing since partitions do not have to be constructed.

Hamlet and Taylor [10] later investigated the same question and concluded that the Duran/Ntafos model was unrealistic. One reason was that the overall failure probability was too high. Hamlet and Taylor theoretically strengthened the case for partition testing but made an important point that partition testing can be no better than the information that defines the subdomains. Gutjahr [11] followed up on Hamlet and Taylor's results and showed theoretically that partition testing is consistently more effective than random testing under realistic assumptions.

More recent results have been produced that favor partition testing over random testing in practical cases. Reid [6] and Yin, Lebne-Dengel, and Malaiya [12] both performed experiments with different partition testing strategies and compared them with random testing. In all cases random testing was found to be less effective

```
1. Identify functional units that may be tested separately.
    For each functional unit, identify parameters and environment variables
    that affect the behavior of the functional unit.
2. Identify choices for each parameter and environment individually.
3. Determine constraints among the choices.
4. Generate all combinations, so called test frames, of parameter choices
    that satisfy the constraints.
5. Transform the test frames into test cases by instantiating the choices.
```

Figure 1: Overview of the steps of the category partition method.
than the investigated partition testing methods.
A key issue in any partition testing approach is how partitions should be identified and how values should be selected from them. In early partition testing methods like Equivalence Partitioning (EP) [2] and Boundary Value Analysis (BVA) [2], parameters of the test problem are identified. Each parameter is then analyzed in isolation to determine suitable partitions of that parameter. Support for identifying parameters and their partitions from arbitrary specifications is rather limited in EP and BVA. Ostrand and Balcer proposed the Category Partition (CP) method partially to address this problem [1].

### 2.1 Category Partition

The category partition method [1] consists of a number of manual steps by which a natural language specification is systematically transformed into an equivalence class model for the parameters of the test object. Figure 1 shows the steps of the category partition method.

Constraints allow the tester to decrease the size of the test suite. Even with constraints, the number of test frames generated by all combinations can be combinatorial $\left(\prod_{i=1}^{N} v_{i}\right.$, where $v_{i}$ is the number of values for each one of the $N$ parameters). Combination strategies can be used to further decrease the number of test cases in the test suite.

## 3 Combination Strategies - Methods

Combination strategies is a class of test case selection methods where test cases are identified by choosing "interesting" values ${ }^{1}$, and then combining those values of test object parameters. The combinations are selected based on some combinatorial strategy. Some combination strategies are based on techniques from experimental design.

This section first explains the different coverage criteria, normally associated with combination strategies and then briefly describes the combination strategies that were identified in the literature. The combination

[^1]

Figure 2: Classification scheme for combination strategies.
strategies have been organized into different classes based on the amount of randomness of the algorithm and according to how the test suites are created. Figure 2 shows an overview of the classification scheme. The combination strategies labeled non-deterministic all depend to some degree on randomness. A property of these combination strategies is that the same input parameter model may lead to different test suites. The simplest non-deterministic combination strategy is pure random selection of test cases. The group of non-deterministic combination strategies also include two heuristic methods, Simulated Annealing (SA) and Automatic Efficient Test Generator (AETG). Finally, two combination strategies are based on Artificial Life (AL) techniques, Genetic Algorithm (GA) and Ant Crawl Algorithm (ACA).

The deterministic combination strategies group is further divided into two subgroups, instant and iterative. All of these combination strategies will always produce the same result from a specific input parameter model.

The two instant combination strategies, Orthogonal Arrays (OA) and Covering Arrays (CA), produce the complete test suite directly.

The iterative combination strategies build the test suite step by step. The parameter-based combination strategy In Parameter Order (IPO) starts by creating a test suite for a subset of the parameters in the input parameter model. Then one parameter at a time is added and the test cases in the test suite are modified to cover the new parameter. Completely new test cases may also need to be added.

The largest group of combination strategies is test case based. They share the property that the algorithms generate one test case at a time and add it to the test suite. Each Choice (EC), Partly Pair-Wise (PPW), Base Choice (BC), All Combinations (AC), Anti-random (AR), CATS, $k$-perim, and $k$-bound all belong to the test case based category.

### 3.1 Coverage Criteria for Combination Strategies

Like many test case selection methods, combination strategies are based on coverage. In the case of combination strategies, coverage is determined with respect to the values of the parameters of the test object that the tester decides are interesting. The simplest coverage criterion, each-used, does not take into account how interesting values of different parameters are combined, while the more complex coverage criteria, such as pair-wise coverage, is concerned with (sub-)combinations of interesting values of different parameters. The following subsections define the coverage criteria satisfied by combination strategies included in this paper.

Each-used (also known as 1-wise) coverage is the simplest coverage criterion. $100 \%$ each-used coverage requires that every interesting value of every parameter is included in at least one test case in the test suite.
$100 \%$ Pair-wise (also known as 2-wise) coverage requires that every possible pair of interesting values of any two parameters are included in some test case. Note that the same test case may cover more than one unique pair of values.

A natural extension of pair-wise (2-wise) coverage is $t$-wise coverage, which requires every possible combination of interesting values of $t$ parameters be included in some test case in the test suite. $t$-wise coverage is formally defined by Williams and Probert [13].

A special case of $t$-wise coverage is $N$-wise coverage, where $N$ is the number of parameters of the test object. $N$-wise coverage requires all possible combinations of all interesting values of the $N$ parameters be included in the test suite.

The each-used, pair-wise, $t$-wise, and $N$-wise coverage criteria are purely combinatorial and do not use any semantic information. However, semantic information may also be used when defining coverage criteria. Cohen et al. [14] indicate that normal and error parameter values should be treated differently with respect to coverage. Normal values lie within the bounds of normal operation of the test object, and error values lie outside of the normal operating range. Often, an error value will result in some kind of error message and the termination of the execution. To avoid one error value masking another Cohen et al. suggest that only one error value of any parameter should be included in each test case. This observation was also made and explained in an experiment by Grindal et al. [15].

A special case of $t$-wise coverage called variable strength was proposed by Cohen, Gibbons, Mugridge, and Colburn [16]. This strategy requires higher coverage among a subset of parameters and a lower coverage criterion across all variables. Assume for example a test problem with four parameters $A, B, C, D$. Variable strength may require 2 -wise coverage across all parameters and 3 -wise coverage for parameters $B, C, D$.

By considering only the valid values, a family of coverage criteria corresponding to the general $t$-wise coverage criteria can be obtained. For instance, $100 \%$ each valid used coverage requires every valid value of every parameter be included in at least one test case in which the rest of the values also are valid. Correspondingly, $100 \% t$-wise valid coverage requires every possible combination of valid values of $t$ parameters be included in some test case, and the rest of the values are valid.

Error values may also be considered when defining coverage criteria. A test suite satisfies single error coverage if each error value of every parameter is included in some test case in which the rest of the values are valid.

A special case of normal values was used by Ammann and Offutt [17] to define base choice coverage. First, a base test case is identified by choosing an "interesting value" of each parameter, such as the most frequently used value. The assumption is that the most frequently used value of each parameter is a normal value. $100 \%$ base choice coverage requires every interesting value of each parameter be included in a test case in which the rest of the values are base values. Further, the test suite must also contain the base test case.

| Parameter |  | Number |  |
| ---: | :---: | :---: | :---: |
|  | $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ |
| $\mathbf{A}$ | 2 | 27 | abc |
| $\mathbf{B}$ | 8 | 15 |  |
| $\mathbf{C}$ | + | - |  |
| $\mathbf{D}$ | do not use | 1 | 3 |

Table 1: Mapping of parameters and parameter values for the series example. An empty cell means no existing value.

### 3.2 The Test Problem

The following sections describe a number of combination strategies. A common example is introduced here, and this example is used to illustrate the various strategies.

Consider a command "series <start> <stop> [<step>]," which generates a sequence of numbers that begin at <start> and end at <stop> with an optional <step>. For example, the command "series 515 3" will generate the sequence "5 811 14."

The first step for the tester is to identify a set of parameters and interesting values to use as input to the combination strategies. Since the intent of this example is to illustrate the strategies, values are selected to be small yet illustrative, thus may not be complete for testing.

The problem is modeled with four parameters $\langle A, B, C, D\rangle$. $A$ is used for different starting values, $B$ is used for different stopping values, $C$ is used for the sign ( $+/-$ ) of the step, and $D$ is used for different step values. Table 1 gives values for each parameter. The abstract test case $\langle 1,2,1,3\rangle$ represents the 1 st value for $A$, the 2nd value for $B$, the 1 st for $C$, and the 3 rd for $D$, thus it represents the command series 215 +3 .

An empty cell, for example $B 3$, means that no parameter value is defined for the testing problem. Whenever that value is included in a test case by a combination strategy it represents a "free choice" and may be exchanged for an arbitrary defined value, in this case $B 1$ or $B 2$.

The "abc" value of $A 3$ represents an invalid value, that is, characters instead of integers. The "do not use" value of $D 1$ means that the step value should be omitted from the command.

The following sections use only the abstract representation of the test problem, that is, A1, B2, etc.

### 3.3 Non-Deterministic Combination Strategies

Non-deterministic combination strategies all share the property that chance will play a certain role in determining which tests are generated. This means two executions of the same algorithm with the same preconditions may produce different results.

### 3.3.1 Heuristic Combination Strategies

## Heuristic t-wise (AETG)

Burroughs, Jain, and Erickson [18] and Cohen, Dalal, Kajla, and Patton [19] report on the use of a tool called Automatic Efficient Test Generator (AETG), which contains an algorithm for generating all pair-wise

```
Assume test cases }\mp@subsup{t}{1}{}-\mp@subsup{t}{i-1}{}\mathrm{ already selected
Let UC be a set of all pairs of values of any two parameters that are
not yet covered by the test cases t1 - ti-1
A) Select candidates for ti by
    1) Selecting the variable and the value included in most
        pairs in UC.
    2) Put the rest of the variables into a random order.
    3) For each variable in the sequence determined by step two,
        select the value included in most pairs in UC.
B) Repeat steps 1-3 k times and let ti be the test case that covers
    the most pairs in UC. Remove those pairs from UC.
Repeat until UC is empty.
```

Figure 3: AETG heuristic algorithm for achieving pair-wise coverage.
combinations. Cohen et al. [20,14] later described a heuristic greedy algorithm for achieving $t$-wise coverage, where $t$ is an arbitrary number. This algorithm was implemented in the AETG tool. Figure 3 shows an algorithm that will generate a test suite to satisfy pair-wise coverage.

There are 36 possible combinations for the example in section 3.2. Suppose that test cases $\{(\mathrm{A} 1, \mathrm{~B} 2, \mathrm{C} 1, \mathrm{D} 2)$, (A2, B2, C1, D3), (A3, B1, C2, D1) \} have already been selected. Then $U C$ contains the set of pairs $\{(\mathrm{A} 1, \mathrm{~B} 1),(\mathrm{A} 1, \mathrm{C} 2),(\mathrm{A} 1, \mathrm{D} 1),(\mathrm{A} 1, \mathrm{D} 3),(\mathrm{A} 2, \mathrm{~B} 1),(\mathrm{A} 2, \mathrm{C} 2),(\mathrm{A} 2, \mathrm{D} 1),(\mathrm{A} 2, \mathrm{D} 2),(\mathrm{A} 3, \mathrm{~B} 2),(\mathrm{A} 3, \mathrm{C} 1)$, (A3, D2), (A3, D3), (B1, C1), (B1, D2), (B1, D3), (B2, C2), (B2, D1), (C1, D1), (C2, D2), (C2, D3) \}.

When selecting candidates for the next test case, parameter values are evaluated to find which is included in the most pairs in $U C . B 1$ and $C 2$ each appear in 5 pairs. Suppose the algorithm selects $C 2$, yielding the partial test case ( $\left.,{ }_{-}, \mathrm{C} 2,_{-}\right)$Then a random order of the remaining parameters is established. Assume the algorithm selects $D, A$, and $B$. Comparing the possible values of $D,(-,, \mathrm{C} 2, \mathrm{D} 1)$ does not cover any pairs in $U C$, while both $(-,, \mathrm{C} 2, \mathrm{D} 2)$ and $(-,, \mathrm{C} 2, \mathrm{D} 3)$ cover one pair. Assume the algorithm selects $D 2$. Next is to compare the values of $A .(\mathrm{A} 1, \ldots, \mathrm{C} 2, \mathrm{D} 2)$ and (A3, , C2, D2) each cover one additional pair in $U C$ (only pairs including the currently studied parameter are counted). ( $\mathrm{A} 2,,_{,}, \mathrm{C} 2, \mathrm{D} 2$ ) on the other hand covers two pairs in $U C$. Thus, the algorithm selects (A2, , C2, D2). Finally, the possible values of $B$ are investigated. (A2, B1, C2, D2) covers two combinations in $U C$, while (A2, B2, C2, D2) only covers one. Thus, the algorithm has found a candidate, which in total covers 5 pairs in $U C$. The steps to identify candidates are repeated $k$ times. Note, that there is some variance in the order of the remaining parameters and when two or more parameter values cover the same number of pairs of $U C$; in these cases an arbitrary choice is made. After generating enough candidates, the one that includes most pairs in $U C$ is selected as the next test case and the search is restarted for the next test case.

The number of test cases generated by the AETG algorithm for a specific test problem is related to the number of constructed candidates ( $k$ in the algorithm in figure 3) for each test case. In general, larger values of $k$ yield smaller numbers of test cases. However, Cohen et al. report that using values larger than 50 will

```
Input: A partial test set
Output: A test case
Begin
    Create an initial population }P\mathrm{ consisting of }m\mathrm{ candidates
    Evaluate P
    While (stopping condition is not met) {
        Identify Elite for survival consisting of \sigma best individuals from P.
        Apply Selection to individuals in P to create P Pmating,
        consisting of ( }m-\sigma)\mathrm{ individuals.
        Crossover }\mp@subsup{P}{\mathrm{ mating}}{
        MutateP Pating
        P=Elite + P Pmating
        Evaluate P.
        If (stagnation condition is met) Massively mutate P.
    }
    Return the best test case found.
End
```

Figure 4: Genetic algorithm for achieving t -wise coverage.
not dramatically decrease the number of test cases.
The AETG algorithm in figure 3 generates a test suite that satisfies pair-wise coverage. It is straightforward to modify the algorithm to achieve $t$-wise coverage.

## Simulated Annealing (SA)

Cohen, Gibbons, Mugridge, and Colburn [16] suggest using Simulated Annealing (SA) to generate test suites for $t$-wise coverage. The algorithm was not described completely enough to include in this survey.

### 3.3.2 Artificial Life Based Combination Strategies

## Genetic Algorithm (GA)

A genetic algorithm was proposed by Shiba, Tsuchiya, and Kikuno [21]. The algorithm was inspired by AETG [20], and adds one test case at a time to a test suite until enough test cases have been identified to satisfy the desired coverage criterion. Figure 4 shows the genetic algorithm used to find the next test case.

A fitness function based on the number of new combinations covered by the test case candidate plays a

| Parameter | Bits | Value | Significance |
| :--- | :--- | :--- | :--- |
| A | $b 1, b 0$ | 00 | A1 (2) |
|  |  | 01 | A1 (2) |
|  |  | 10 | A2 (27) |
|  |  | 11 | A3 (abc) |
| B | $b 2$ | 0 | B1 (8) |
|  |  | 1 | B2 (15) |
| C | $b 3$ | 0 | C1 (+) |
|  |  | 1 | C2 (-) |
| D | $b 5, b 4$ | 00 | D1 (do not use) |
|  |  | 01 | D2 (1) |
|  |  | 10 | D3 (3) |
|  |  | 11 | D3 (3) |

Figure 5: Encoding of the test problem from section 3.2 into a bit vector.
crucial role in this algorithm. The fitness function is used both to identify the $\sigma$ best individuals for survival to the next iteration (Elite) and in the selection mechanism to identify a reproduction base $P_{\text {mating }} . P_{\text {mating }}$ is created by repeatedly selecting two test cases from $P$, and letting the fitter one win with some probability. Note that the same test case may be selected and included in $P_{\text {mating }}$ several times. The contents of $P_{\text {mating }}$ is then manipulated by uniform crossover and mutation. Crossover is achieved by letting two test case candidates from $P_{\text {mating }}$ exchange values of each position independently with probability 0.5 . Mutation replaces the value of a position in a test case with a randomly chosen value.

Should several iterations not result in any further improvements of the current solution, there is a possibility to "massively mutate" $P$ to avoid getting stuck at a local optimum. The algorithm terminates when the stopping condition is met, in this case, when $M$ test case candidates have been generated. The best candidate among these is then returned.

To illustrate the genetic algorithm, consider the mapping of the test problem in section 3.2 presented in figure 5. Assume that test cases $\{(110100)$, (010010), (101101), (101010) \} have already been selected. Further assume that the population in the generic algorithm is of size 4, and that two test cases are selected to survive in each iteration.

Thus, four candidate test cases are generated randomly in the first step of the algorithm. Assume these candidates are $\{100101,110000,010100,111010\}$. Each candidates is then evaluated using the fitness function (how many new pairs are covered by that candidate). In order, the fitness of the four candidates are $\{0,1,2,0\}$. Thus, the Elite is $\{110000,010100\}$. $P_{\text {mating }}$ is then constructed in two steps. First a pair of candidates is selected randomly from $P$. Suppose the first two candidates are 100101 and 110000. A weighted choice favoring the best fit, 110000 , is then made between the two, and suppose it is selected. Assume in the second iteration that the first and third candidates are selected. The weighted choice favors the third candidate, but assume for the sake of illustration that the first candidate, 100101, is selected. Thus $P_{\text {mating }}$ initially consists of 110000 and 100101. In the crossover, b0 and b4 are selected for crossover, resulting in an updated $P_{\text {mating }}$ of 100001 and 110100 . In the mutation step, b1 of the first candidate is selected for mutation, resulting in the final $P_{\text {mating }}$ of 100011 and 110100. The new content of $P$ is then


Figure 6: The test problem in section 3.2 represented as a graph for the ant colony algorithm.
$\{100011,110000,010100,110100\}$, with the fitness $\{3,1,2,0\}$. In this manner the algorithm iterates until a stopping criterion is reached and the best test case candidate, the one with the highest fitness, is upgraded to a final test case and the algorithm starts over to find a new test case.

GA can be used to construct test suites to satisfy $t$-wise coverage for arbitrary $t$.

## Ant Colony Algorithm (ACA)

The ant colony algorithm was proposed by Shiba, Tsuchiya, and Kikuno [21]. It is based on an algorithm originally designed to solve the traveling salesman problem, and was inspired by the AETG algorithm. One test case at a time is added to a test suite until enough test cases have been identified to satisfy the desired coverage criterion.

The test case is represented as a graph with nodes representing the different parameters in the test problem and edges representing the different values of each parameter. Figure 6 shows the graph that represents the test problem described in section 3.2. A test case is represented as a path from the starting node to the goal node.

The main idea of the algorithm is to let ants crawl the graph, guided by the amount of pheromone at each edge in the graph. When an ant has completed a path, the corresponding test case candidate is evaluated and more pheromone is deposited, along the way, proportional to the quality of the candidate. The complete algorithm can be seen in figure 7. In the early iterations of the algorithm, not enough pheromone has been laid, therefore local heuristics are computed to guide the ants in these iterations.

To illustrate the algorithm, consider the example in figure 6. Assume that test cases \{(A1, B2, C1, D3), (A2, B1, C1, D2), (A1, B2, C2, D3), (A2, B1, C2, D3) \} have already been selected, and local heuristics are based on the inverse of how many times each edge is already used in some test case. Thus, the initial pheromone is laid according to the following: $\{\mathrm{A} 1: 2, \mathrm{~A} 2: 2, \mathrm{~A} 3: 4, \mathrm{~B} 1: 2, \mathrm{~B} 2: 2, \mathrm{C} 1: 2, \mathrm{C} 2: 2, \mathrm{D} 1: 4, \mathrm{D} 2: 2, \mathrm{D} 3: 1, \mathrm{D} 4: 4\}$. Also assume that only one ant at a time crawls the graph.

Assume that the ant selects the path $\{\mathrm{A} 3, \mathrm{~B} 1, \mathrm{C} 2, \mathrm{D} 1\}$, based on the initial pheromone weights. This candidate covers five uncovered pairs, which is a good result so a lot of pheromone is deposited on these edges. The resulting pheromone trail then is $\{\mathrm{A} 1: 2$, $\mathrm{A} 2: 2, \mathrm{~A} 3: 9, \mathrm{~B} 1: 7, \mathrm{~B} 2: 2, \mathrm{C} 1: 2, \mathrm{C} 2: 7, \mathrm{D} 1: 9, \mathrm{D} 2: 2, \mathrm{D} 3: 1, \mathrm{D} 4: 4\}$. The next ant has a higher probability of crawling along the same path as the first. However, there is always an element of chance. Suppose that the next ant selects path $\{\mathrm{A} 3, \mathrm{~B} 1, \mathrm{C} 2, \mathrm{D} 4\}$. This candidate also covers five uncovered pairs, thus a lot of pheromone is again deposited on the trail, resulting in $\{\mathrm{A} 1: 2, \mathrm{~A} 2: 2, \mathrm{~A} 3: 14, \mathrm{~B} 1: 12, \mathrm{~B} 2: 2, \mathrm{C} 1: 2, \mathrm{C} 2: 12, \mathrm{D} 1: 9, \mathrm{D} 2: 2, \mathrm{D} 3: 1, \mathrm{D} 4: 9\}$. When the total number of candidates generated exceeds a predetermined number the algorithm stops and returns the best candidate.

```
Input: A partial test set
Output: A test case
Begin
    Compute local heuristics.
    Initialize pheromone.
    While (stopping condition is not met) {
        For (each ant k) {
            Generate a candidate test case S}\mp@subsup{S}{k}{
            Evaluate }\mp@subsup{S}{k}{
            Lay pheromone
        }
        Update pheromone
        If (stagnation condition is met) Initialize pheromone.
    }
    Return the best test case found.
End
```

Figure 7: Ant colony algorithm for achieving t-wise coverage.

ACA can be used to construct test suites to satisfy $t$-wise coverage for arbitrary $t$.

### 3.3.3 Random Combination Strategies

## Random (Rand)

Creating a test suite by randomly sampling test cases from the complete set of test cases based on some input distribution (often uniform distribution) is an old idea with unclear origins. In the scope of combination strategies, Mandl [22] may have been the first to mention the idea. Mandl states (without motivation) that a reasonable random selection would probably be about the same size as a test suite that simply takes the $N$ variables in turn, and includes for each of them $v$ test cases to cover the $v$ interesting values of the variable.

### 3.4 Deterministic Combination Strategies

Deterministic combination strategies share the property that they produce the same test suite every time they are used. These strategies are divided into two subcategories: instant, iterative combination strategies. Instant combination strategies create the complete test suite all at once. Iterative combination strategies build up the test suite iteratively. The test case based combination strategies add one test case at a time while the parameter-based combination strategies build up the test suite by adding values to the test cases for one parameter at a time.

```
Assume a test problem with N parameters, P}\mp@subsup{P}{1}{},\mp@subsup{P}{2}{},\ldots,\mp@subsup{P}{N}{}
Let }\mp@subsup{P}{1}{}\mathrm{ and }\mp@subsup{P}{2}{}\mathrm{ represent the two most significant parameters.
A) Create all combinations of values of }\mp@subsup{P}{1}{}\mathrm{ and }\mp@subsup{P}{2}{
B) For each parameter }\mp@subsup{P}{3}{}\ldots\mp@subsup{P}{N}{
    Extend the combinations obtained in step A with values of the current parameter
    If combinations remain when all parameter values have been used,
    reuse parameter values until all combinations have been extended.
    If there are not enough combinations for each parameter value to be used
    Reuse existing combinations until all parameter values have been used.
```

Figure 8: An instance of an algorithm implementing the partly pair-wise algorithm.

### 3.4.1 Iterative Combination Strategies

Iterative combination strategies are those in which the test suite is built up gradually.

## Iterative Combination Strategies

In the test case based iterative combination strategies one test case at a time is generated and added to the test suite. Thus, a tester may start the algorithm with an already preselected set of test cases.

## Each Choice (EC)

The basic idea behind the Each Choice combination strategy is to include each value of each parameter in at least one test case. This is achieved by generating test cases by successively selecting unused values for each parameter. This strategy was invented by Ammann and Offutt [17], who also suggested that EC can result in an undesirable test suite since the test cases are constructed mechanically without considering any semantic information.

Applying EC to the example described in section 3.2 would yield the test cases $\{(\mathrm{A} 1, \mathrm{~B} 1, \mathrm{C} 1, \mathrm{D} 1)$, (A2, B2, C2, D2), (A3, B1, C1, D3) \}, where the values of $B$ and $C$ in the third test case is selected arbitrarily. It is clear from the definition of EC that it satisfies each-used coverage.

## Partly Pair-Wise (PPW)

Burroughs, Jain, and Erickson [18] sketch a "traditional" combination strategy in which all pair-wise combinations of the values of the two most significant parameters should be included, while at the same time including each value of the other parameters at least once. No details of the actual algorithm are given, thus the algorithm in 8 is a guess.

Applying the PPW algorithm in figure 8 to the example in section 3.2 yields the test suite $\{(\mathrm{A} 1, \mathrm{~B} 1, \mathrm{C} 1, \mathrm{D} 1)$, (A1, B2, C2, D2), (A2, B1, C1, D3), (A2, B2, C1, D1), (A3, B1, C1, D1), (A3, B2, C1, D1) \}, assuming that A and B are the two most significant parameters.

PPW satisfies variable strength coverage combining pair-wise and each used coverage.

## Base Choice (BC)

The Base Choice (BC) combination strategy was proposed by Ammann and Offutt [17]. The first step of BC is to identify a base test case. The base test case combines the most "important" value for each parameter. Importance may be based on any predefined criterion such as most common, simplest, smallest, or first. Ammann and Offutt suggest "most likely used." Similarly, Cohen et al. called these values "default values." From the base test case, new test cases are created by varying the values of one parameter at a time while keeping the values of the other parameters fixed on the values in the base test case.

A test suite generated by BC satisfies each-used coverage since each value of every parameter is included in the test suite. If "most likely used" is used to identify the base test case, BC also satisfies the single error coverage criterion. The reason is that the base test case only consists of normal values and only one parameter at a time differs from the base test case.

Applying BC to the example in section 3.2 requires a base test case to be identified; (A1, B2, C2, D3). The entire test suite can be derived from the base test case: $\{(\mathrm{A} 1, \mathrm{~B} 2, \mathrm{C} 2, \mathrm{D} 3)$, (A2, B2, C2, D3), (A3, B2, C2, D3), (A1, B1, C2, D3), (A1, B2, C1, D3), (A1, B2, C2, D1), (A1, B2, C2, D2) \}

Ammann and Offutt state that satisfying base-choice coverage does not guarantee the adequacy of the test suite for a particular application. However, a strong argument can be made that a test suite that does not satisfy base-choice coverage is inadequate.

A slight variation of the base choice combination strategy was later described by Cohen et al. [19]. In their version, called "default testing," the tester varies the values of one parameter at a time while the other parameters contain some default value.

The term "default testing" has also been used by Burr and Young [23], describing yet another variation of the general base choice theme. In their version all parameters except one contains the default value, and the remaining parameters contain a maximum or a minimum value. This variant does not necessarily satisfy each-used coverage.

## All Combinations (AC)

The All Combinations (AC) combination strategy algorithm generates every combination of values of the different parameters. The origin of this method is impossible to trace back to a specific person due to its simplicity. It is often used as a benchmark with respect to the number of test cases $[22,17,14]$.

An AC test suite satisfies $N$-wise coverage.

## Antirandom (AR)

Antirandom testing was proposed by Malaiya [24], and is based on the idea that test cases should be selected to have maximum "distance" from each other. Parameters and interesting values of the test object are encoded using a binary vector such that each interesting value of every parameter is represented by one or more binary values. Test cases are then selected such that a new test case resides on maximum Hamming or Cartesian distance from all the other already selected test cases.

To illustrate the concept of AR and the difference between Hamming and Cartesian distance consider the test problem described in section 3.2. The first step is to encode the values of the test object into a binary vector. Parameters $B$ and $C$ have two variables each, thus one single bit per parameter is sufficient for the encoding. Parameters $A$ and $D$ have three variables each. For these parameters at least two bits per parameter are needed. Using two bits for a parameter provides four different values. Thus, the tester may choose to add another value or let two values of the two-bit sequence represent the same parameter value. Table 9 shows an example of a complete mapping, in which one value of parameter A is used twice and a new value for parameter D has been added to illustrate both concepts.

Suppose the bit vector (001010), corresponding to (A2, B1, C2, D1), has been selected as the starting

| Parameter | Bits | Value | Significance |
| :--- | :--- | :--- | :--- |
| A | $b 1, b 0$ | 00 | A1 (2) |
|  |  | 01 | A1 (2) |
|  |  | 10 | A2 (27) |
|  |  | 11 | A3 (abc) |
| B | $b 2$ | 0 | B1 (8) |
|  |  | 1 | B2 (15) |
| C | $b 3$ | 0 | C1 (+) |
|  |  | 1 | C2 (-) |
| D | $b 5, b 4$ | 00 | D1 (do not use) |
|  |  | 01 | D2 (1) |
|  |  | 10 | D3 (3) |
|  |  | 11 | $-(5)$ |

Figure 9: A second encoding of the test problem from section 3.2 into a bit vector.
point for AR in the example. The AR algorithm is now invoked to select the second test case. Either the Hamming distance or the Cartesian distance are used to find a test case as far away as possible from the starting test case. The Hamming distance is the number of bits in which two binary vectors differ. The Cartesian distance is defined as the square root of the Hamming distance. In the example, it is straightforward to see that the second bit vector selected by the AR is the bit-inverse, (110101). The Hamming distance between the two vectors is 6 and the Cartesian distance is $\sqrt{6}$.

The third bit vector is selected such that the sum of the distances to the first two is maximized. If Hamming distance is used, the sum of the distances to the two selected bit vectors is 6 for any bit vector, i.e., any bit vector may be selected as the third. The reason is that each bit in the third bit vector differs from exactly one of the already selected bit vectors. Thus, the sum of the distances is either $1+5,2+4,3+3,4+2$, or $5+1$. If, on the other hand, Cartesian distance is used, only bit vectors with exactly half of the bits coming from each already selected vector are candidates to be selected as the third bit vector. The reason is that $\sqrt{3}+\sqrt{3}$ is larger than both $\sqrt{2}+\sqrt{4}$ and $\sqrt{1}+\sqrt{5}$.

When selecting the fourth bit vector, the sum of the distances to all of the three previously selected vectors is considered and so on.

In the original description, the antirandom testing scheme is used to create a total order among all the possible test cases rather than attempting to limit the number of test cases. However, just like it is possible in random testing to set a limit to how many test cases that should be generated, the same applies to antirandom test cases. Thus, antirandom testing can be used to select a subset of all possible test cases, while ensuring that they are as far apart as possible.

In the general case, a test suite generated by AR does not satisfy any of the normal coverage criteria associated with combination strategies.

## k-boundary (k-bound) and k-perimeter (k-perim)

With a single parameter restricted to values $[L, . ., U]$ (ordered from L (ower) to $\mathrm{U}($ pper )), generating boundary values is easy. The tester can select the set $\{L, U\}$ or $\{L, L+1, \ldots, U-1, U\}$ or $\{L, L+1, L+2, \ldots, U-$
$2, U-1, U\}$ etc. With multiple domains two problems arise: (1) the number of points grows rapidly and (2) there are several possible interpretations. Hoffman, Strooper and White [25] defined the k-boundary and k -perimeter combination strategies, where k is an arbitrary integer, to handle these problems.

The 1-boundary of an individual domain is the smallest and largest values of that domain ( $L$ and $U$ ). The 2-boundary of an individual domain is the next to the smallest and the next to the largest values of that domain $(L+1$ and $U-1)$ and so on.

The $k$-boundary of a set of domains is the Cartesian product of the k-boundaries of the individual domains. Thus, the number of test cases of the k-boundary is $2^{N}$, where $N$ is the number of domains. In the general case, k-boundary does not satisfy any of the normal coverage criteria associated with combination strategies.

The $k$-perimeter test suite can be constructed by (1) including the k-boundary test suite; (2) forming all possible pairs of test cases from the k-boundary test suite such that the two test cases in a pair only differ in one dimension, e.g., $(L, L, L, L)$ and $(U, L, L, L)$; and (3) for each identified pair of test cases, adding to the k-perimeter test suite all points in between the two test cases of the pair i.e., $\{(L+1, L, L, L),(L+$ $2, L, L, L), \ldots,(U-2, L, L, L),(U-1, L, L, L)\}$.

The example used to illustrate the majority of the other combination strategies is not well suited to illustrate k -bound and k -perim due to the small number of parameter values for some of the parameters. Thus, consider a test problem with two parameters $A$ and $B$. Suppose $A$ has seven parameter values $\{\mathrm{A} 1, \mathrm{~A} 2$, A3, A4, A5, A6, A7\} and $B$ has four parameter values $\{\mathrm{B} 1, \mathrm{~B} 2, \mathrm{~B} 3, \mathrm{~B} 4\}$. Also assume that the parameter values are ordered from smallest to largest. The following test suites satisfy $100 \% 1$-bound and 2 -bound coverage respectively: $\{(\mathrm{A} 1, \mathrm{~B} 1),(\mathrm{A} 1, \mathrm{~B} 4),(\mathrm{A} 7, \mathrm{~B} 4),(\mathrm{A} 7, \mathrm{~B} 1)\}$ and $\{(\mathrm{A} 2, \mathrm{~B} 2),(\mathrm{A} 2, \mathrm{~B} 3),(\mathrm{A} 6, \mathrm{~B} 3),(\mathrm{A} 6, \mathrm{~B} 2)\}$. 3 -bound and higher is not relevant in this example since parameter $B$ only contains four values.

Adding all the intermediate test cases to the test suite satisfying $100 \%$ 1-bound coverage yield $100 \% 1$ perim coverage: $\{(\mathrm{A} 1, \mathrm{~B} 1),(\mathrm{A} 1, \mathrm{~B} 2)$, (A1, B3), (A1, B4), (A2, B4), (A3, B4), (A4, B4), (A5, B4), (A6, B4), (A7, B4), (A7, B3), (A7, B2), (A7, B1), (A6, B1), (A5, B1), (A4, B1), (A3, B1), (A2, B1)\}. In the same manner a test suite for $100 \%$ 2-perim coverage can be created: $\{(\mathrm{A} 2, \mathrm{~B} 2),(\mathrm{A} 2, \mathrm{~B} 3),(\mathrm{A} 3, \mathrm{~B} 3),(\mathrm{A} 4, \mathrm{~B} 3)$, (A5, B3), (A6, B3), (A6, B2), (A5, B2), (A4, B2), (A3, B2) \}
k -perimeter does not satisfy any coverage criterion in the general case, however, 1-perimeter satisfies each-used coverage.

## Heuristic t-wise (CATS)

The CATS tool for generating test cases is based on a heuristic algorithm that can be custom designed to satisfy $t$-wise coverage. The algorithm was described by Sherwood [26] and a version of the algorithm that generates a test suite to satisfy pair-wise coverage is shown in figure 10 .

There are 36 possible combinations for the example in section 3.2. Suppose that test cases $\{(\mathrm{A} 1, \mathrm{~B} 2, \mathrm{C} 1, \mathrm{D} 2)$, (A2, B2, C1, D3), (A3, B1, C2, D1) \} have already been selected. Then $Q$ contains the remaining 33 combinations. $U C$ contains the set of pairs $\{(\mathrm{A} 1, \mathrm{~B} 1),(\mathrm{A} 1, \mathrm{C} 2),(\mathrm{A} 1, \mathrm{D} 1),(\mathrm{A} 1, \mathrm{D} 3),(\mathrm{A} 2, \mathrm{~B} 1),(\mathrm{A} 2, \mathrm{C} 2)$, (A2, D1), (A2, D2), (A3, B2), (A3, C1), (A3, D2), (A3, D3), (B1, C1), (B1, D2), (B1, D3), (B2, C2), (B2, D1), (C1, D1), (C2, D2), (C2, D3) \}.

Among the combinations in $Q$ there are two, (A1, B1, C2, D3) and (A2, B1, C2, D2), that contain 5 pairs still in $U C$, which is the current maximum. Thus, the algorithm selects (A1, B1, C2, D3) as the next test case and it is removed from $Q$. The new contents of $U C$ is then $\{(\mathrm{A} 1, \mathrm{D} 1),(\mathrm{A} 2, \mathrm{~B} 1),(\mathrm{A} 2, \mathrm{C} 2),(\mathrm{A} 2, \mathrm{D} 1)$, (A2, D2), (A3, B2), (A3, C1), (A3, D2), (A3, D3), (B1, C1), (B1, D2), (B2, C2), (B2, D1), (C1, D1), $(\mathrm{C} 2, \mathrm{D} 2)\}$. In the next iteration, $(A 2, B 1, C 2, D 2)$ is the only test case in $U C$ that contains 5 pairs. Thus, it is selected and the resulting contents of $U C$ is now $\{(\mathrm{A} 1, \mathrm{D} 1),(\mathrm{A} 2, \mathrm{D} 1),(\mathrm{A} 3, \mathrm{~B} 2),(\mathrm{A} 3, \mathrm{C} 1),(\mathrm{A} 3, \mathrm{D} 2),($ $\mathrm{A} 3, \mathrm{D} 3),(\mathrm{B} 1, \mathrm{C} 1),(\mathrm{B} 2, \mathrm{C} 2),(\mathrm{B} 2, \mathrm{D} 1),(\mathrm{C} 1, \mathrm{D} 1)\}$. Of the remaining 31 combinations in $Q$, only one still in

```
Assume test cases }\mp@subsup{t}{1}{}-\mp@subsup{t}{i-1}{}\mathrm{ already selected
Let Q be the set of all possible combinations not yet selected
Let UC be a set of all pairs of values of any two parameters that are
not yet covered by the test cases t}\mp@subsup{t}{1}{}-\mp@subsup{t}{i-1}{
A) Select ti by finding the combination that covers most pairs in UC.
    If more than one combination covers the same amount select the
        first one encountered.
    Remove the selected combination from Q .
    Remove the covered pairs from UC.
B) Repeat until UC is empty.
```

Figure 10: CATS heuristic algorithm for achieving pair-wise coverage.
$U C,(\mathrm{~A} 3, \mathrm{~B} 2, \mathrm{C} 1, \mathrm{D} 1)$, contains 4 pairs. The process repeats until $U C$ is empty.
The resulting set of test cases in this example is $\{(\mathrm{A} 1, \mathrm{~B} 2, \mathrm{C} 1, \mathrm{D} 2),(\mathrm{A} 2, \mathrm{~B} 2, \mathrm{C} 1, \mathrm{D} 3),(\mathrm{A} 3, \mathrm{~B} 1, \mathrm{C} 2, \mathrm{D} 1)$, (A1, B1, C2, D3), (A2, B1, C2, D2), (A3, B2, C1, D1), (A1, B1, C1, D1), (A2, B2, C2, D1), (A3, B1, C1, D2), (A3, B1, C1, D3) \}.

The CATS algorithm has some similarities with AETG. The main difference is in how the next test case to be included in the final test suite is found. CATS examines the whole list of unused test cases to find one that adds as much new coverage as possible while AETG constructs test case candidates one parameter value at a time based on coverage information. The constructed candidates are then evaluated to find the best possible, which is added to the test suite.

In CATS it is guaranteed that "best" test case is always selected while in AETG there are no guarantees. However, for large test problems, AETG is more efficient since only a small set of test cases has to be evaluated in each step.

The greedy nature of the algorithm makes it impossible to exactly calculate the number of test cases in a test suite generated by the algorithm.

## Parameter-Based Combination Strategies

There is only one parameter-based combination strategy, In Parameter Order (IPO). Just as with test case based combination strategies, a tester may start the IPO algorithm with an already preselected set of test cases.

## In Parameter Order (IPO)

For a system with two or more parameters, the in-parameter-order (IPO) combination strategy [27, 28, 29] generates a test suite that satisfies pair-wise coverage for the values of the first two parameters. The test suite is then extended to satisfy pair-wise coverage for the values of the first three parameters, and continues to do so for the values of each additional parameter until all parameters are included in the test suite.

To extend the test suite with the values of the next parameter, the IPO strategy uses two algorithms. The first algorithm, horizontal growth was shown in figure 11. It extends the existing test cases in the test

```
Algorithm IPO_H (\tau, pi)
{
    Let \tau be a test suite that satisfies pair-wise coverage for the values of
    parameters p}\mp@subsup{p}{1}{}\mathrm{ to }\mp@subsup{p}{i-1}{}\mathrm{ .
    Assume that parameter pi contains the values }\mp@subsup{v}{1}{},\mp@subsup{v}{2}{},\ldots,\mp@subsup{v}{q}{
    \pi={ pairs between values of pi and values of }\mp@subsup{p}{1}{}\mathrm{ to }\mp@subsup{p}{i-1}{}
    if (| \tau | \leqq)
    {
        for 1\leqj\leq|\tau|, extend the jth test in \tau by adding value vj}\mathrm{ and
        remove from }\pi\mathrm{ pairs covered by the extended test.
    }
    else
    {
        for 1\leqj\leqq, extend the jth test in \tau by adding value vj}\mathrm{ and
        remove from \pi pairs covered by the extended test
        for q<j\leq| \tau|, extend the jth test in }\tau\mathrm{ by adding one value of pi
        such that the resulting test covers the most number of pairs in }\pi\mathrm{ ,
        and remove from }\pi\mathrm{ pairs covered by the extended test
    }
}
```

Figure 11: IPO_H - An algorithm for horizontal growth of a test suite by adding values for new parameters.
suite with values of the next parameter. The second algorithm, vertical growth, shown in figure 12, creates additional test cases such that the test suite satisfies pair-wise coverage for the values of the new parameter.

To illustrate how the IPO algorithm works, consider the test problem described in section 3.2. The starting point for IPO is a test suite that satisfies pair-wise coverage for a subset of the parameters of the test problem. Consider parameters $A$ and $B$. A test suite that satisfies pair-wise coverage for these parameters is $\{(\mathrm{A} 1, \mathrm{~B} 1),(\mathrm{A} 1, \mathrm{~B} 2),(\mathrm{A} 2, \mathrm{~B} 1),(\mathrm{A} 2, \mathrm{~B} 2),(\mathrm{A} 3, \mathrm{~B} 1),(\mathrm{A} 3, \mathrm{~B} 2)\}$. Then IPO_H is invoked to extend the existing test cases with the values of parameter $C$.

When invoking IPO_H, $\tau$ contains the test suite and $\pi$ contains all uncovered pairs including the values of $C,\{(\mathrm{~A} 1, \mathrm{C} 1),(\mathrm{A} 1, \mathrm{C} 2),(\mathrm{A} 2, \mathrm{C} 1),(\mathrm{A} 2, \mathrm{C} 2),(\mathrm{A} 3, \mathrm{C} 1),(\mathrm{A} 3, \mathrm{C} 2),(\mathrm{B} 1, \mathrm{C} 1),(\mathrm{B} 1, \mathrm{C} 2),(\mathrm{B} 2, \mathrm{C} 1),(\mathrm{B} 2, \mathrm{C} 2)\}$. The first step is to add the first value of $C$ to the first test case in $\tau$, the second value of C to the second test case in $\tau$, etc., until each value of $C$ has been used once. Thus, $\tau=\{(\mathrm{A} 1, \mathrm{~B} 1, \mathrm{C} 1)$, (A1, B2, C2), (A2, B1, $)$, $\left.(\mathrm{A} 2, \mathrm{~B} 2,)_{-}\right),(\mathrm{A} 3, \mathrm{~B} 1,)_{)}\left(\mathrm{A} 3, \mathrm{~B} 2,,_{)}\right)$. Covered pairs are removed from $\pi$ resulting in $\{(\mathrm{A} 2, \mathrm{C} 1),(\mathrm{A} 2, \mathrm{C} 2)$, (A3, C1), (A3, C2), (B1, C2), (B2, C1) \}.

The next step of IPO_H is to extend the remaining test cases in $\tau$ with a value of $C$ to cover as many pairs as possible in $\pi$. Thus, for the third test case, (A2, B1, C1) covers one pair in $\tau$ while (A2, B1, C2) covers two, hence the latter is selected and the covered pairs are removed from $\pi$. Repeating these steps for all test cases in $\tau$ yields $\tau=\{(\mathrm{A} 1, \mathrm{~B} 1, \mathrm{C} 1),(\mathrm{A} 1, \mathrm{~B} 2, \mathrm{C} 2),(\mathrm{A} 2, \mathrm{~B} 1, \mathrm{C} 2),(\mathrm{A} 2, \mathrm{~B} 2, \mathrm{C} 1),(\mathrm{A} 3, \mathrm{~B} 1, \mathrm{C} 1)$,

```
Algorithm IPO_V (\tau,\pi)
{
    Where }\tau\mathrm{ contains the set of already selected test cases
    and }\pi\mathrm{ contains the set of still uncovered pairs
    Let }\tau\prime\mathrm{ be an empty set
    for each pair in }
    { assume that the pair contains value w of pk, 1\leqk<i, and value u of pi
    {
        if ( }\tau\mathrm{ ' contains a test case with ' '-') as the value of p
        and }u\mathrm{ as the value of }\mp@subsup{p}{i}{}\mathrm{ )
            modify this test case by replacing the ''-') with w
    else
            add a new test case to }\tau\prime\mathrm{ that has }w\mathrm{ as the value of }\mp@subsup{p}{k}{},u\mathrm{ as
            the value of pi, and ''_') as the value of every other parameter;
    }
    \tau=\tau\cup\tau\prime
}
```

Figure 12: IPO_V - An algorithm for vertical growth of a test suite by adding values for new parameters.
(A3, B2, C2) $\}$ and $\pi=\{ \}$. Since $\pi$ is empty, $\tau$ satisfies pair-wise coverage also for the values of $C$, which means that IPO_V does not need to be invoked.

In the next iteration, parameter $D$ is considered. Again, IPO_H is invoked, this time with $\tau=\{(\mathrm{A} 1, \mathrm{~B} 1, \mathrm{C} 1)$, (A1, B2, C2), (A2, B1, C2), (A2, B2, C1), (A3, B1, C1), (A3, B2, C2) \} and $\pi=\{(\mathrm{A} 1, \mathrm{D} 1),(\mathrm{A} 1, \mathrm{D} 2)$, (A1, D3), (A2, D1), (A2, D2), (A2, D3), (A3, D1), (A3, D2), (A3, D3), (B1, D1), (B1, D2), (B1, D3), (B2, D1), (B2, D2), (B2, D3), (C1, D1), (C1, D2), (C1, D3), (C2, D1), (C2, D2), (C2, D3)\}.

Adding the three values of D to the three first test cases of $\tau$ results in $\tau=\{(\mathrm{A} 1, \mathrm{~B} 1, \mathrm{C} 1, \mathrm{D} 1)$, (A1, B2, C2, D2), (A2, B1, C2, D3), (A2, B2, C1, $),(\mathrm{A} 3, \mathrm{~B} 1, \mathrm{C} 1, ~),,(\mathrm{A} 3, \mathrm{~B} 2, \mathrm{C} 2, ~-)\}$ and $\pi=\{(\mathrm{A} 1, \mathrm{D} 3)$, (A2, D1), (A2, D2), (A3, D1), (A3, D2), (A3, D3), (B1, D2), (B2, D1), (B2, D3), (C1, D2), (C1, D3), (C2, D1) \}. Extending the remaining three test cases of $\tau$ with values of $D$ such that as many pairs in $\pi$ are covered and removed results in $\tau=\{(\mathrm{A} 1, \mathrm{~B} 1, \mathrm{C} 1, \mathrm{D} 1),(\mathrm{A} 1, \mathrm{~B} 2, \mathrm{C} 2, \mathrm{D} 2),(\mathrm{A} 2, \mathrm{~B} 1, \mathrm{C} 2, \mathrm{D} 3),(\mathrm{A} 2, \mathrm{~B} 2, \mathrm{C} 1, \mathrm{D} 1)$, (A3, B1, C1, D2), (A3, B2, C2, D1) $\}$ and $\pi=\{(\mathrm{A} 2, \mathrm{D} 2),(\mathrm{A} 3, \mathrm{D} 1),(\mathrm{A} 3, \mathrm{D} 3),(\mathrm{B} 2, \mathrm{D} 3),(\mathrm{C} 1, \mathrm{D} 3)\}$.

Since $\pi$ is non-empty there are still some pairs involving parameter values of $D$ that are not covered by the test suite. Thus, IPO_V is invoked with $\tau$ and $\pi$ as above. IPO_V iterates over the pairs of $\pi$. In the first iteration, $\tau \prime$ is empty, which results in a new partial test case being added to $\tau \prime$ such that the current pair of $\pi$ is covered, (A2, , , D2). The next pair of $\pi,(\mathrm{A} 3, \mathrm{D} 1)$ is incompatible with the first test case in $\tau \prime$, thus a new incomplete test case is added, with the result that $\tau \prime=\left\{\left(\mathrm{A} 2,_{-},{ }_{-}, \mathrm{D} 2\right),\left(\mathrm{A} 3,_{-}, \quad \mathrm{D} 1\right)\right\}$. In the third iteration, $\tau \prime$ is further extended with (A3, , , D3). In the fourth iteration, the current pair (B2, D3) can be fitted into the third test case of $\tau \prime$, changing it to (A3, B2, , D3). Finally, the last pair of $\pi$ is incompatible with the existing three, resulting in a fourth incomplete test case being added to $\tau \prime$, with the result $\left\{\left(\mathrm{A} 2,_{-},-\mathrm{D} 2\right),\left(\mathrm{A} 3,_{-}, \mathrm{D} 1\right),\left(\mathrm{A} 3, \mathrm{~B} 2,_{-}, \mathrm{D} 3\right),\left(_{-}, \mathrm{C} 1, \mathrm{D} 3\right)\right\}$.

| 1 | 2 | 3 |
| :--- | :--- | :--- |
| 2 | 3 | 1 |
| 3 | 1 | 2 |

Figure 13: A $3 \times 3$ latin square.

| 1 | 2 | 3 |
| :--- | :--- | :--- |
| 2 | 3 | 1 |
| 3 | 1 | 2 |$\quad$| 1 | 2 | 3 |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 3 | 1 | 2 |
| 2 | 3 | 1 | | 1,1 | 2,2 | 3,3 |
| :--- | :--- | :--- |
| 2,3 | 3,1 | 1,2 |
| 3,2 | 1,3 | 2,1 |

Figure 14: Two orthogonal $3 \times 3$ latin squares and the resulting combined square.

In the last step of the IPO algorithm the test cases of $\tau \prime$ is added to the already existing test cases of $\tau$, resulting in a test suite that satisfies pair-wise coverage for all four parameters. The undecided parts of the incomplete test cases in the final test suite need to be instantiated at execution time, but in terms of pair-wise coverage any value can be used, so it is up to the tester.

The nature of the IPO algorithms makes it difficult to theoretically calculate the number of test cases in a test suite generated by the algorithm.

An algorithm closely resembling IPO was informally described by Huller [30].

### 3.4.2 Instant Combination Strategies

Sometimes a tester wants to include her favorite test cases in the test suite. With instant combination strategies this will not affect the final result. The same test suite will be selected by the combination strategy regardless if some other test cases have already been preselected.

## Orthogonal Arrays (OA)

Orthogonal Arrays is a mathematical concept that has been known for quite some time. The application of orthogonal arrays to testing was first introduced by Mandl [22] and later more thoroughly described by Williams and Probert [31].

The foundation of OA is Latin Squares. A Latin Square is a $V \times V$ matrix completely filled with symbols from a set that has cardinality $V$. The matrix has the property that the same symbol occurs exactly once in each row and column. Figure 13 contains an example of a $3 \times 3$ Latin Square with the symbols $\{1,2,3\}$.

Two Latin Squares are orthogonal if, when they are combined entry by entry, each pair of elements occurs precisely once in the combined square. Figure 14 shows an example of two orthogonal $3 \times 3$ Latin Squares and the resulting combined square.

If indexes are added to the $\mathrm{r}(\mathrm{ows})$ and the c (olumns) of the matrix, each position in the matrix can be described as a tuple $\left\langle r, c, z_{i}\right\rangle$, where $z_{i}$ represents the values of the $\langle r, c\rangle$ position. Figure 15 contains the indexed Latin Square from Figure 14 and Figure 16 contains the resulting set of tuples. The set of all tuples constructed by a Latin Square satisfies pair-wise coverage.

The set of test cases in figure 16 can be used directly for the example described in section 3.2 by mapping the dimensions of the tuple $<r c z_{1} z_{2}>$ on the parameters $A, B, C$, and $D$. Recall that parameters $A$ and $D$ contain three values each, while parameters $B$ and $C$ only contain two parameters each. Assuming a mapping that preserves the stated order, that is, that maps $A$ onto $r, B$ onto $c, C$ onto $z_{1}$, and $D$ onto $z_{2}$,

| Coordinates | c(olumns) |  |  |
| :---: | :---: | :---: | :---: |
| $\mathbf{r}(\mathbf{o w s )}$ | $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ |
| $\mathbf{1}$ | 1,1 | 2,2 | 3,3 |
| $\mathbf{2}$ | 2,3 | 3,1 | 1,2 |
| $\mathbf{3}$ | 3,2 | 1,3 | 2,1 |

Figure 15: Two orthogonal $3 \times 3$ latin squares augmented with coordinates.

| tuple | Tuple $<r c z_{1} z_{2}>$ |
| :---: | :---: |
| $\mathbf{1}$ | 1111 |
| $\mathbf{2}$ | 1222 |
| $\mathbf{3}$ | 1333 |
| $\mathbf{4}$ | 2123 |
| $\mathbf{5}$ | 2231 |
| $\mathbf{6}$ | 2312 |
| $\mathbf{7}$ | 3132 |
| $\mathbf{8}$ | 3213 |
| $\mathbf{9}$ | 3321 |

Figure 16: Tuples from the two orthogonal $3 \times 3$ latin squares that satisfies pair-wise coverage.
test cases $3,5,6,7,9$ all contain coordinate values without corresponding parameter values. This problem is resolved simply by substituting the undefined coordinate value with an arbitrary defined value. Consider for example test case 7 in Figure 16. The value of coordinate $z_{1}$ is 3 while only values 1 and 2 are defined in the mapping from parameter $C$. To create a test case from tuple 7 , the undefined value should be replaced by an arbitrarily defined value, which in this case is 1 or 2 .

A test suite based on orthogonal arrays satisfies pair-wise coverage, even after undefined values have been replaced and possibly some duplicate tuples have been removed. Williams and Probert [31] give further details on how test cases are created from orthogonal arrays.

Sloane has collected a large number of precalculated orthogonal arrays of various sizes and made them available on the $\mathrm{Web}^{2}$.

## Covering Arrays (CA)

Covering Arrays [32] is an extension of orthogonal arrays. A property of orthogonal arrays is that they are balanced, which means that each parameter value occurs the same number of times in the test suite. If only $t$-wise (for instance pair-wise) coverage is desired the balance property is unnecessary and will make the algorithm less efficient. In a covering array that satisfies $t$-wise coverage, each $t$-tuple occurs at least once but not necessarily the same number of times. Another problem with orthogonal arrays is that for some problem sizes there do not exist enough orthogonal arrays to represent the entire problem. This problem is also avoided by using covering arrays.

A covering array is denoted $O(r, c, k)$, where $r$ is the number of rows (test cases), $c$ is the number of columns (parameters), and $k$ is the maximum values of any parameter. A covering array is constructed for a test problem from a set of building blocks:

- $O(r, c, k)$ is an orthogonal array based test suite with $r$ rows (test cases) and $c$ columns (parameters), which have at most $k$ values each. In addition, for this orthogonal array to be useful, $r, c$, and $k$ should be related by the formulas $r=k^{2}$ and $c=k+1$.
- $B(r-1, c, k, d)=O(r, c, k)$ with the first row removed and with the columns used $d$ times each consecutively.
- $R(r-k, c-1, k, d)=O(r, c, k)$ with the $k$ first rows and the first column removed and with the remaining columns used $d$ times each consecutively.
- $I(r, c)$ is a matrix with $r$ rows and $c$ columns completely filled with "ones."
- $N(r, k, c)$ is a matrix with $r$ rows and $c$ columns, which of which consists of $k-1$ submatrices of size $k$ rows $\times c$ columns filled with "twos", "threes", etc. up to $k$.
Figure 17 gives examples of the five types of building blocks.
Depending on the size of the problem, these five building blocks are combined in different ways to create a test suite with pair-wise coverage over all parameters and parameter values. Figure 18 shows some examples of how building blocks may be combined. The algorithm to identify and assemble suitable building blocks for a given problem size is omitted from this presentation for reasons of size but may be found Williams' paper [32]. Both the algorithms for finding suitable building blocks and for combining these to fit a specific test problem have been automated.

An extension of the work on CA was given by Williams and Probert [33]. A similar approach also building on the concept of building blocks to construct a covering array was proposed by Kobayashi, Tsuchiya, and Kikuno [34].

[^2]

Figure 17: Examples of the five types of building blocks for generation of covering arrays.


Figure 18: Examples of how different building blocks may be combined to covering arrays. $O$ is short for $O(9,4,3)$.


Figure 19: Subsumption hierarchy for the coverage criteria related to combination strategies. Combination strategies that support a certain coverage criterion is indicated by italics.

## 4 Comparing Combination Strategies

This section contains comparisons of combination strategies with respect to coverage criteria and size of generated test suite

### 4.1 Subsumption

A subsumption hierarchy for the combination coverage criteria reported in the literature is shown in Figure 19. The definition of subsumption is from Rapps and Weyuker [35]: coverage criterion X subsumes coverage criterion Y iff $100 \%$ X coverage implies $100 \%$ Y coverage (Rapps and Weyukers' original paper used the term inclusion instead of subsumption). A guiding principle for this work has been to generalize the known coverage criteria to obtain a complete picture.

The left, boxed-in, column of coverage criteria in Figure 19 represent criteria that do not use semantic information. $N$-wise coverage requires that every combination of all parameter values are included in the test suite. " $t$-wise coverage" is used as a short form for every level of coverage from $N-1$ down to 2 . It is straightforward to see that $j$-wise coverage subsumes $(j-1)$-wise coverage, for any $j$ such that $2 \leq j \leq N$.

The right column contains the coverage criteria based only on valid values. For the same reason that $j$-wise coverage subsumes $(j-1)$-wise coverage, it is easy to see that the $j$-wise valid coverage criterion subsumes the $(j-1)$-wise valid coverage criterion.

Note that $i$-wise coverage where $1 \leq i \leq N-1$ does not subsume $i$-wise valid coverage. This is easily demonstrated for the case where $i=2$. Consider a three parameter problem, each with one valid value $v$ and one erroneous value $e$. The following four combinations satisfy pair-wise coverage since each possible pair of values is included $[(v, v, e),(v, e, v),(e, v, v),(e, e, e)]$. However, pair-wise valid coverage is not satisfied since a pair of two valid parameters only counts if the remaining parameters are also valid. Since every combination in the example contains at least one erroneous value, the achieved pair-wise valid coverage is $0 \%$. $N$-wise coverage contains every possible combination so by definition it subsumes all other coverage criteria, including $N$-wise valid coverage.

The single error coverage criterion represents a different class of coverage criteria. Although it is possible to generalize this idea into multiple error coverage, this would introduce the possibility of masking as was described in section 3.1, so it is omitted. At first it may seem that the single error coverage criterion would subsume the each-used coverage criterion. However, a counter example can be constructed. Consider a twoparameter problem where the first parameter has two valid values (v11 and v12) and one erroneous value (e11). The second parameter has one valid value (v21) and one erroneous value (e21). The two test cases $[(e 11, v 21),(v 11, e 21)]$ satisfy the single error coverage criterion but not the each-used coverage criterion.

An interesting property of base choice coverage, when assuming that all base choices are valid, is that it subsumes the each-used, each valid used, and single error coverage criteria. By definition, the base test case contains only valid values. From the definition of the base choice algorithm it follows that all test cases differ from the base test case by the value of only one parameter. Further, each interesting value of every parameter is included in some test case. Thus, the each-used criterion is subsumed. Further, every valid value of every parameter must appear in at least one test case in which the rest of the values are also valid, since the base test case contains only valid values. Thus, the each valid used criterion is subsumed. Finally, for the same reason, every invalid value of every parameter must occur in exactly one test case in which the rest of the values are valid. Thus, the single error criterion is subsumed.

### 4.2 Size of Generated Test Suite

Table 2 gives an overview of the sizes of the generated test suites. In some cases only approximate values are possible to give. The main reason for this is that those algorithms contain some degree of randomness. Thus, the same combination strategy may produce different solutions to the same test problem. Also, only a subset of the described combination strategies is included in the table. There are several reasons. Some strategies, i.e., Rand and AR, have no natural limit. Other strategies, i.e., k-bound and k-perim, differ greatly in the size of the generated test suite depending on the value of k . The algorithm of PPW is not described and finally CATS, having an element of greediness, is not investigated well enough to be described in terms of size of generated test suite.

In addition to the formulas relating the size of the test problem to the size of the generated test suite, table 2 contains two example test problems one with few parameters and many values and one with many parameters with few values.

| Combination <br> Strategy | Test Suite <br> Size | Example 1 <br> $N=8, V_{i}=4$ | Example 2 <br> $N=4, V_{i}=8$ |
| :--- | :--- | ---: | ---: |
| EC | $V_{\max }$ | 4 | 8 |
| BC | $1+\sum_{i=1}^{N}\left(V_{i}-1\right)$ | 25 | 29 |
| pair-wise (IPO, <br> AETG, OA, CA <br> GA, ACA, SA) | $\sim V_{\max }^{2}$ | $\sim 16$ | $\sim 64$ |
| AC | $\prod_{i=1}^{N} V_{i}$ | 65536 | 4096 |

Table 2: Definite or approximate ( $\sim$ ) test suite sizes for a problem with $N$ parameters and $V_{i}$ values of the $i$ :th parameter. $V_{\max }$ is the largest $V_{i}$.

## 5 Combination Strategies - Experience and Evaluations

This section gives a brief overview of the collected experience of using combination strategies. First some theoretical results are described, then a brief overview of implemented tools is given. The most important real-life applications of combination strategies are then presented and finally some results from comparing different combination strategies are given.

### 5.1 Applicability of Combination Strategies - Theory

Heller [36] uses a realistic example to show that testing all combinations of parameter values is infeasible in practice. Heller concludes that there is a need to identify a subset of combinations of manageable size. The author suggests fractional factorial designs, which are closely related to the orthogonal arrays combination strategies. Traditionally, these have been used in natural science experiments to decrease the number of experiments. The underlying mathematics of the fractional factorial designs allows researchers to establish which factors cause which results. It is also shown that the number of test cases from fractional factorial designs is less than all combinations.

Dalal and Mallows [37] provide a theoretical view of combination strategies. They give a model for software faults in which faults are classified according to how many parameters (factors) need distinct values to cause the fault to result in a failure. A $t$-factor fault is triggered whenever the values of some $t$ parameters are involved in triggering a failure. Each possible fault can be specified by giving the combination(s) of values relevant for triggering that fault.

Following the lead of Dalal and Mallows, Kuhn and Reilly [38] gives a strong argument for the use of combination strategies. They investigated 365 error reports from two large real life projects and discovered that pair-wise coverage was nearly as effective in finding faults as testing all combinations, since most faults were 1 - or 2 -factor faults. Their findings were strengthened by Kuhn and Wallace [39] in another study of the relationship between input and output values in a number of real applications.

Dunietz, Ehrlich, Szablak, Mallows, and Iannino [40] examined the correlation between $t$-wise coverage and achieved code coverage. A number of test suites for a specific test problem were created by random sampling. For every test suite, $t$-wise coverage with $t=1,2, \ldots, n$ were calculated. Then the test suites were executed and block (statement) and path coverage were computed. The results of this study is that

```
Assume a test problem with three parameters a, b, and c, each
with 2 possible values.
Let tc1 = (a1, b1, c1), tc2=(a2, b2, c1), and tc3=(a1, b1, c2)
The following nine pairs are covered by the three test cases:
(a1,b1), (a1,c1), (b1,c1), (a2,b2), (a2,c1), (b2,c1), (a1,b1), (a1,c2), (b1,c2)
Note that the pair (a1,b1) occurs both in tc1 and tc3.
The following four pairs are not covered by the three test cases:
(a1,b2), (a2,b1), (a2,c2), (b2,c2)
2-coverage = 8/12=67%
2-diversity = 8/9 = 89%
```

Figure 20: An example of 2-coverage and 2-diversity of a test suite.
test suites satisfying $t$-wise coverage, where $t \geq 2$, all perform similarly terms of block (statement) coverage. Thus, the authors are providing an argument for using pair-wise coverage instead of all-combination coverage.

Piwowarski, Ohba, and Caruso [41] describe how to successfully apply code coverage as a stopping criterion during functional testing. The authors formulated functional testing as a problem of selecting test cases from a set of possible test cases made up of all combinations of values of the input parameters. Burr and Young [23] show that continually monitoring achieved code coverage works as a tool to improve the input parameter model. Initial experiments showed that ad hoc testing resulted in about $50 \%$ block and decision coverage. By continually applying code coverage as the input parameter models were refined, decision coverage was increased to $84 \%$.

Dalal, Jain, Karunanithi, Leaton, Lott, Patton, and Horowitz [42] present an architecture of a generic test-generation system based on combination strategies. A test data model is derived from the requirements and then used to generate test input tuples. Test cases, containing test inputs and expected results are generated by the combination strategy from the test data model. The test inputs are fed to the system under test, which generates actual output. Finally the actual and expected outputs are compared.

The authors claim that the model-based testing depends on three key technologies: the notation used for the data model, the test generation algorithm, and the tools for generating the test execution infrastructure including expected output.

Dalal and Mallows [37] provide two measures of efficiency of a combination strategy test suite. The $t$-wise coverage (in Dalal's and Mallow's words $t$-coverage) is the ratio of distinct $t$-parameter combinations covered to the total number of possible $t$-parameter combinations.

The $t$-diversity is the ratio of distinct $t$-parameter combinations covered to the total number of $t$-parameter combinations in the test suite. Thus, this metric summarizes to what degree the test suite avoids replication. Figure 20 shows an example of these two metrics.

Williams and Probert [13] formally define a set of coverage criteria, including $t$-wise coverage, through
the use of a so called interaction element, which may be used in the same way statements and decisions are used for code coverage. Williams and Probert show that the general problem of finding a minimal set of test cases that satisfies $t$-wise coverage can be NP-complete.

In the seminal work on antirandom testing, Malaiya [24] observes that for some combination strategies, the order in which test cases are generated may help preserve execution. In the case of antirandom testing, the algorithm always finds the next test case as far away as possible from the already existing test cases. If that order is preserved during execution, and if the execution is prematurely terminated, the executed test cases will always be optimally spread across the input space.

### 5.2 Tools for Combination Strategies

Tools have been implemented for some combination strategies, although only one is available commercially. T-GEN [43] is a test case generator implemented for the Category Partition method (CP) [1]. T-GEN may also be used to generate test cases based on the output space in a similar manner as the input space is handled in CP. Also, T-GEN can generate complete or partial test scripts under certain conditions.

The Orthogonal Arrays Testing System (OATS) tool [44] is based on the OA combination strategy. OATS automatically identifies orthogonal arrays and generates test suites that satisfy pair-wise coverage of the values of the parameter values.

PairTest [28] is a tool implemented in Java. It implements the IPO combination strategy for pair-wise coverage.

Telcordia (Formerly Bellcore) has a test tool called Automatic Efficient Test Generator (AETG) [45]. It implements the heuristic combination strategy. This tool has been used for both research and practice since AETG is also commercially available on the web ${ }^{3}$.

Daley, Hoffman, and Strooper [46] describe the "Roast" framework. This approach differs from the previously described tools by attempting to automate both the generation and execution of test cases. The Roast framework consists of four steps. (1) Generate creates the combinations of parameter values to be tested, that is, the input portions of each test case. (2) Filter removes invalid combinations from the test suite. (3) Execute drives the execution of the test cases. (4) Check compares the expected with the actual output. Apart from supplying a parameter model of the test problem, the only manual intervention is adding the expected output.

### 5.3 Applicability of Combination Strategies - Practice

Plenty of examples of the applicability of combination strategies in software testing can be found in the literature. Within the scope of functional testing, Dalal, Jain, Karunanithi, Leaton, Lott, Patton and Horowitz [47, 42] report results from using the AETG tool. The tool was used to generate test cases for Bellcore's Intelligent Service Control Point, a rule-based system used to assign work requests to technicians, and a GUI window in a large application. In a previous study, Cohen, Dalal, Kajla and Patton [19] demonstrated the use of AETG for screen testing, by testing the input fields for consistency and validity across a number of screens. Dalal et al. [47] report four lessons learned. First, the model of the test data is fundamental and requires considerable domain expertise. Second, model-based testing is in fact a development project since it consists of much testware that needs to be developed and maintained. Third, change must be handled and

[^3]unnecessary manual intervention should be avoided. Fourth, technology transfer requires careful planning, that is, changing a work process or a state-of-practice is a project in itself.

Burr and Young [23] also used the AETG tool. They tested a third party email product from Nortel that converts email messages from one format to another. Another example of using combination strategies for functional testing is described by Huller [30], who used an IPO related algorithm to test ground systems for satellite communications.

Closely related to functional testing is the use of the k-bound and k-perim combination strategies within the Roast Framework [46] to semi-automatically test Java Classes.

Although the examples of using combination strategies for functional testing dominate, there are several examples of the applicability of combination strategies in other areas of testing. Williams and Probert [31] demonstrate how combination strategies can be used to select configurations for configuration testing. The authors point out that the cost for setting up a certain configuration is often substantial. Further, each configuration requires complete test suites to be executed, which adds to the cost of testing the configuration. Thus, it is important to find ways to identify a small number of configurations to test. In this case OA was successfully used. A related attempt is by Yilmaz, Cohen and Porter [48]. They use covering arrays as a starting point for fault localization in complex configuration spaces. When a fault of unclear origin occurs, covering arrays are used to select a set of configurations. The same test cases are then executed on all configurations and the results are analyzed, to localize the origin of the fault.

Kropp, Koopman, and Siewiorek [49] describe the Ballista testing methodology, which supports automated robustness testing of off-the-shelf software components. Robustness is defined as the degree to which a software component functions correctly in the presence of exceptional inputs or stressful conditions. Ballista generates a test suite that has all combinations of all parameter values. If the total number of combinations is too large to be feasible, 5000 in the reported experiment, a deterministic pseudo-random sample is used. The application of the Ballista methodology is demonstrated on several implementations of the POSIX operating system C language API.

Berling and Runesson [50] describe how to use fractional factorial designs for performance testing. Specifically, they focus on how to find a minimal or near minimal set of test cases to determine the effect on the system performance of different parameter values. An important observation by Berling and Runesson is that when the result variable has a nominal type, (for example, true or false, or pass or fail as in testing) the property of fractional factorial designs cannot be exploited.

### 5.4 Comparisons of Combination Strategies

Several papers report results from comparing combination strategies. This section contains brief descriptions and conclusions of these studies.

Huller [30] shows that pair-wise configuration testing may save more than $60 \%$ in both cost and time compared to quasi-exhaustive testing. A similar conclusion was reached by Brownlie, Prowse, and Phadke [44] who compared the results of using OA on one version of a PMX/StarMAIL release with the results from conventional testing on a prior release. The authors estimated that $22 \%$ more faults would have been found if OA had been used on the first version, despite the fact that conventional testing required twice as much time to perform.

Several studies compare the performance of different combination strategies. By far the most popular property used to compare combination strategies is number of test cases generated for a specific test problem. In particular, this is interesting for the non-deterministic and greedy combination strategies, AETG, SA, GA, and ACA, since the size of the test suite may not be determined algebraically. A number of papers
compare a subset of the combination strategies satisfying 2 -wise or 3 -wise coverage: IPO and AETG [27], OA and AETG [15], CA and IPO [32], and AETG, IPO, SA, GA, and ACA [21, 16]. Interesting to note is that in all these comparisons, these combination strategies perform similarly with respect to the number of test cases. SA seems to be generating the least number of test cases most of the time, but the differences are usually small.

Since the performance with respect to generated test cases does not largely favor one particular combination strategy, several of the authors have also compared the strategies with respect to the execution time for the combination strategy to generate its test cases. Lei and Tai [28] show that the time complexity of IPO is superior to the time complexity of AETG. IPO has a time complexity of $O\left(v^{3} N^{2} \log (N)\right)$ and AETG has a time complexity of $O\left(v^{4} N^{2} \log (N)\right)$, where $N$ is the number of parameters, each of which has $v$ values. Further, Williams [32] reports that CA outperforms IPO by almost three orders of magnitude for the largest test problems in their study, in terms of time taken to generate the test suites. Finally, Shiba et al. [21] show some execution times, but the executions have been made on different target machines so the results are a bit inconclusive. However, SA, SGA and ACA seems to perform similarly, more than an order of magnitude faster that AETG.

Grindal et al. [15] take another approach to compare combination strategies, the number of faults found. They show that BC performs as well, in terms of detecting faults, as AETG and OA despite fewer test cases for a number test objects seeded with 131 faults. However, similar fault detection ratio does not necessarily mean that the combination strategies detect the same faults. The authors show that BC and OA and AETG target different types of faults.

Yet another way to evaluate combination strategies is on the basis of achieved code coverage of the generated test suites. This approach was taken by Cohen et al. [20]. They report on test suites generated by AETG for 2-wise coverage that reach over $90 \%$ block coverage. Burr and Young [23] reached similar results for AETG in another study. AETG reached $93 \%$ block coverage with 47 test cases, compared to $85 \%$ block coverage for a restricted version of BC using 72 test cases.

Yin, Lebne-Dengel, and Malaiya [12] also used code coverage to compare AR with randomly selected test cases. The antirandom combination strategy consistently performed well by reaching high coverage with few test cases.

## 6 Conclusions and Future Directions

This survey is an attempt to collect and describe in a comprehensive manner some of the properties of combination strategies for test generation. In particular, this survey gives an overview of the coverage criteria normally associated with combination strategies and shows how these coverage criteria are interrelated in a subsumption hierarchy. This survey also relates the surveyed combination strategies based on the approximate sizes of the test suites generated by the combination strategies. It reports on the usefulness of combination strategies and it also suggests a classification scheme for the existing combination strategies based on properties of their algorithms.

Despite the many papers covering combination strategies, several important issues still remain largely unexplored. One such issue is the identification of suitable parameters and parameter values. At a first glance, identifying the parameters of a test problem seems like an easy task. Almost all software components have some input parameters, which could be used directly. This is also the case in most of the papers on combination strategies [49]. However, Yin, Lebne-Dengel, and Malaiya [12] point out that in choosing a set of parameters the problem space should be divided into sub-domains that conceptually can be seen as
consisting of orthogonal dimensions. These dimensions do not necessarily map one-to-one onto the actual input parameters of the implementation. Along the same lines of thinking, Cohen, Dalal, Parelius, and Patton [20] state that in choosing the parameters, one should model the system's functionality, not its interface. It is an open question how the effectiveness and efficiency of the testing is affected by the choice of parameters and their values.

A second issue relating to the use of combination strategies not adequately investigated is constraints among the parameter values. It is not unusual in a test problem that certain subcombinations of parameter values are not allowed. For some of the combination strategies, for example AETG [14], it is possible to extend the algorithms to avoid selecting combinations that are not allowed. It is easy to see that this approach is not feasible for the instant combination strategies, for example OA [31]. Another approach is to rewrite the test problem into several conflict free sub problems [14, 31, 46]. A third approach is to change one or more of the values involved in an invalid combination [17]. These and possible other suggestions to handle constraints in the test problem have not been adequately investigated. Important questions are which constraint handling methods work for which combination strategies and how the size of the test suite is affected by the constraint handling method.

A third issue that needs further investigation is how to select an appropriate combination strategy for a particular test problem. As shown in section 5.4, several comparisons have been made. Some general advice on the choice of combination strategies has been given. For instance, Offutt and Ammann [17] state that the each-choice (EC) combination strategy may result in less effective test suites since test cases are assembled without considering any semantic information, thus the test cases actually selected by the combination strategy may have little to do with how the application is actually used in operation.

Beizer [3] explained that test cases that contain multiple invalid parameter values are normally not effective since most applications will find one of the values and then revert to some error handling routine, ignoring the values of any other parameters. Thus, a fault in the error handling routine for the second invalid value may be masked. Grindal et al. [15] observed that OA and AETG produced this behavior and contributed this to the use of pair-wise strategies over the complete input parameter space. They suggested the combined use of a pair-wise strategy and BC, which ensures that each invalid value of any parameter will be included in at least one test case, where all other parameters are valid, that is, free of masking. Another approach for the same problem is proposed by Cohen et al. [19]. They suggest that the pair-wise combination strategies should be used only for the valid values of each parameter. For invalid values, some other combination strategy ensuring freeness of masking should be used. Although, undoubtedly useful for the tester, these types of information only provides partial advice to the tester when deciding upon which combination strategy to use.

Also some experimental results give advice to the practicing tester. One such advice is that all combination strategies that satisfy pair-wise coverage generate test suites of similar size. Despite these advice, which combination strategies to select is still largely unanswered. In order to make combination strategies widely accessible to testers, future research needs to investigate all of these issues more thoroughly.

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[^1]:    ${ }^{1}$ The term "interesting" may seem insufficiently precise and also a little judgmental. However, it appears frequently in the literature, so is also used in this paper. There are many ways to decide which values are "interesting" and some of the combination strategies discuss criteria for making those choices. In this paper, "interesting values" are whatever values the tester decides to use.

[^2]:    ${ }^{2}$ URL: http://www.research.att.com/~njas/oadir, page last visited November 2004.

[^3]:    ${ }^{3}$ http://aetgweb.agreenhouse.com, page last visited November 2004.

